The Hidden Performance Price of Virtual Functions

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 - Better algorithms
 - Better exploiting the underlying hardware
 - Better usage of the standard library
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- Virtual functions in C++
 - Enable flexibility
 - The basic component of OOP
- Virtual functions are slower than regular functions
- The performance price of virtual functions depends on several factors
 - Here we explain what are those factors









- C++ standard doesn't mandate implementation of virtual functions
- Most compilers, however, implement virtual functions in a similar manner









How virtual functions work - virtual tables









How virtual functions work - function calls











- Virtual functions are more expensive than non-virtual functions
 - The virtual function's address is not known at compile time
 - The program needs to look up the virtual function's address at runtime
 - Virtual function's address lookup is done through virtual table pointer









- A vector of 20 million objects of the same type
- 20 million calls to the virtual function vs 20 million calls to the non-virtual function

	Virtual function call	Non-virtual function call
Short and fast function	153 ms	126 ms
Long and slow function	32.090 ms	31.848 ms









Initial Analysis - conclusion

- The results don't look that bad
- There is a noticeable overhead for small function (18%).
- For the large function, the overhead is negligible
- But is this all there is to virtual functions?









- To activate virtual function mechanism, you need to access the object through a pointer or a reference
 - Objects need to be allocated on the heap (using new, malloc or smart pointers)
- Accessing objects on the heap can be very slow









Optimal layout



Non-optimal layout





- Accessing objects on the heap can be very slow
 - The reason are data cache misses
 - If objects are neighbors in memory, we can expect performance improvements
 - If objects are not neighbors in memory, we can expect slowdowns
 - If the neighboring pointers do not point to neighboring elements on the heap, we can expect data cache misses
- There is no guarantee that the neighboring pointers will point to neighboring objects in memory
 - As the program becomes bigger and more complex there is less and less chance that this will happen
- Vector of objects is much better for the performance compared to vector of pointers
 - The vector of objects doesn't suffer from data cache misses





Vector of pointers - experiment

- Vector of objects containing 20 million objects
- Another vector of pointers, pointer at location *i* points to an object at location *i*
 - \circ $\,$ $\,$ This is the perfect ordering: neighboring pointers point to neighboring objects $\,$
- We measure the time needed to iterate through 20 million objects by following the pointers in the vector of pointers
 - There are several iterations of the experiment
 - \circ $\$ In each new iteration of the experiment, we shuffle the pointer vector a bit
 - Shuffling slows down the traversal a bit
 - We measure the runtime as a function of number of shuffles
 - One shuffle means swap pointer at position [0] with a pointer at position [rand(0, vector_len)]







How swapping of pointers in an array influences the speed of access







Vector of pointers - conclusion

- Memory layout is very important for program performance
 - \circ \quad Worst case is 7.5 times slower than the fastest case
- The slowdown isn't related to virtual functions per se
 - The slowdown is related to the memory layout
 - Still, the main reason you want to use the vector of pointers to achieve polymorphism
- Alternatives to vector of pointers:
 - Use `std::variant` with `std::visitor`
 - Use polymorphic_vector uses virtual dispatching, but doesn't uses pointers. Downside is increased memory consumption → google `polymorphic_vector`
 - Use per type vector (e.g. `boost::base_collection`), a very useful if you don't need a specific ordering in the vector







- Compiler knows the address of non-virtual functions at compile time.
 - This means the compiler can inline the non-virtual function and avoid the function call
- Inlining saves a few instructions on the function call, but that is not all
- After inlining, the compiler can perform many other compiler optimizations, e.g:
 - Move loop invariant code outside of the calling loop
 - Use special instructions that can process more than one data at a time in a process called vectorization <- this can increase speed from 2 to 6 times.









Compiler Optimizations - Example

```
void my loop(int arr len) {
                                                                         void my loop inlined(int arr len) {
   std::vector<double> in;
                                                                            std::vector<double> in;
   std::vector<double> out;
                                                                            std::vector<double> out;
   ...
                                                                            ...
   for (int i = 0; i < out.size(); i++) {</pre>
                                                                            if (debug) {
       out[i] = my sqrt(in[i]);
                                                                                for (int i = 0; i < out.size(); i++) {</pre>
                                                                                     if (in[i] < 0) {
   }
                                                                                         std::cerr << "Value " << a <<</pre>
                                                                         "negative\n";
                                            Loop unswitching
                                                                                    out[i] = std::sqrt(in[i]);
double my sqrt(double a) {
   if (debug && a < 0) {
                                                                              else ·
       std::cerr << "Value " << a << "negative\n";</pre>
                                                                                for (int i = 0; i < out.size(); i++) {</pre>
   }
                                                                                    out[i] = std::sqrt(in[i]);
   return std::sqrt(a);
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```



Compiler Optimizations - experiment

```
class object {
  protected:
    bool m_is_visible;
    unsigned int m_id;
    static unsigned int m_offset;
  public:
    ATTRNOINLINE
    bool is_visible() { return m_is_visible; }
    ATTRNOINLINE
    unsigned int get_id3() { return m_id + m_offset; };
};
```

```
// Test loop
for (int i = 0; i < arr_len; i++) {
    object* o = pv.get(i);
    if (o->is_visible()) {
        count += o->get_id3();
    }
}
```







Compiler Optimizations - results

• Measured the performance of non-virtual function, inlined and non-inlined.

```
// Test loop
for (int i = 0; i < arr_len; i++) {
    object* o = pv.get(i);
    if (o->is_visible()) {
        count += o->get_id3();
    }
}
```

Vector 20M objects	Non-inlined	Inlined
Runtime	242 ms	136 ms







Compiler Optimizations - conclusion

- Virtual functions inhibit compiler optimizations because they are essentially not inlinable
- A solution for this is *type based processing*
 - Don't mix the types, each type has its own container and its own loop
 - The compiler can inline small functions and perform the compiler optimizations
 - Already implemented in `boost::base_collection`
 - This approach is applicable if objects in the vector don't have to be sorted
- The benefits of compiler optimization that happen due to inlining are very case dependent
 - Some code profits a lot from compiler optimizations, other not so much
 - Smaller functions in principle benefit more





Jump Destination Guessing

- To speed up computation, modern CPUs do a lot of guessing (technical term is *speculative execution*)
- In the case of virtual function:
 - The CPU guesses which virtual function will get called
 - \circ $\hfill It starts executing the instructions belonging to the guessed virtual function$
- If the guess is correct, this saves time
- If the guess is wrong, the CPU needs to cancel the effect of wrongly executed instructions and start over
 - This costs time







Jump Destination Guessing - experiment

- Three vector of 20 million objects
 - First vector is sorted by type: A, A, A, A, B, B, B, B, C, C, C, C, D, D, D, D
 - Types in vector in predictable fashion: A, B, C, D, A, B, C, D, A, B, C, D, A, B, C, D
 - Types in vector random: B, C, A, C, A, C, B, B, A, C, B, A.
- We measure time needed to call a small virtual function on the three types of vectors









Performance of virtual functions depending on the sorting type







Jump Destination Guessing - conclusion

- Types sorted in predictable manner -> the CPU can successfully predict the address of the virtual function and this speeds up the computation
- If types are appear randomly, the CPU cannot guess successfully and precious cycles are lost
 - A solution to this is again, type based processing
 - However, type based processing is not always usable
- The effect is mostly pronounced with short virtual functions









Instruction Cache Evictions

- Modern CPUs rely on "getting to know" the instructions they are executing
- The code that has already been executed is hot code
 - Its instructions are in the instruction cache
 - Its branch predictors know the outcome of the branch (true/false)
 - Its jump predictors know the target of the jump
- The CPU is faster when executing hot code compared to executing cold code
- The CPU's memory is limited
 - The code that is currently hot will eventually become cold unless executed frequently
- Virtual functions, especially large virtual functions where each object has a different virtual function, mean that we are switching from one implementation to another
 - The CPU is constantly switching between different implementations and is always running cold code





Instruction Cache Evictions - experiment

- Measuring the effect of instruction cache eviction is the hardest, because it depends on many factors
 - The number of different virtual function implementations the bigger the number, the slower the code
 - The number of executed instructions in the virtual functions the bigger the number, the slower the code
 - The size of virtual function correlates to the number of executed instructions, but they are not the same
 - How sorted are the objects in the container (by type)
 - Best case is when they are sorted by type (AAABBBCCCDDD)
 - Worst case is when they are sorted by type in a round robin fashion (ABCDABCDABCD)







Instruction Cache Evictions - experiment

- Four classes: *rectangle, circle, line* and *monster*
- Four implementations of *long_virtual_functions*
- The *long_virtual_function* consists of a *for* loop with a large *if/elseif/.../else* inside it
- For measurements we use two vectors (20 million objects)
 - Elements of the vector sorted by type: AAABBBCCCDDD
 - Elements of the vector sorted by type in a round-robin fashion:ABCDABCDABCD
- We change the number of comparisons in a large *if/elseif/.../else* block and compare the time needed to iterate the two vectors









Relative performance difference between the sorted and round-robin vector



In the worst case, the same function took 7.5 seconds to execute in the sorted vector, and 12.3 seconds to execute in the round-robin vector





Instruction Cache Eviction - conclusion

- In our example, the cold code was running at the speed of 0.6 of the speed of the fast code
- The phenomenon is not related to the virtual functions themselves
 - E.g, it will happen if each instance has a pointer to a different function
- However, it is most likely to occur with large virtual functions on mixed-type unsorted vectors with many different derived types









- Virtual functions do not incur too much additional cost by themselves
- It is the environment where they run which determines their speed
- The hardware craves predictability: same type, same function, neighboring virtual address
 - When this is true, the hardware run at its fastest
 - It's difficult to achieve this with casual usage of virtual functions
- In game development, they use another paradigm instead of OOP called: data-oriented design
 - One of its major parts is **type based processing**: each vector holds one type only
 - This eliminates all the problems related to virtual functions
 - However, this approach is not applicable everywhere







- If you need to use virtual functions, bear in mind:
 - The number one factor that is responsible for bad performance are data cache misses
 - Avoiding vector of pointers on a hot path is a must!
 - \circ $\,$ $\,$ Other factors also play their role, but to a lesser extend
 - With careful design, you can reap most benefit of virtual functions without incurring too much additional cost
- Here are a few ideas to fix your code with virtual functions:
 - Arrangement of objects in memory is very important!
 - Try to make small functions non-virtual!
 - Most overhead of virtual functions comes from small functions, they cost more to call than to execute
 - Try to keep objects in the vector sorted by type







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