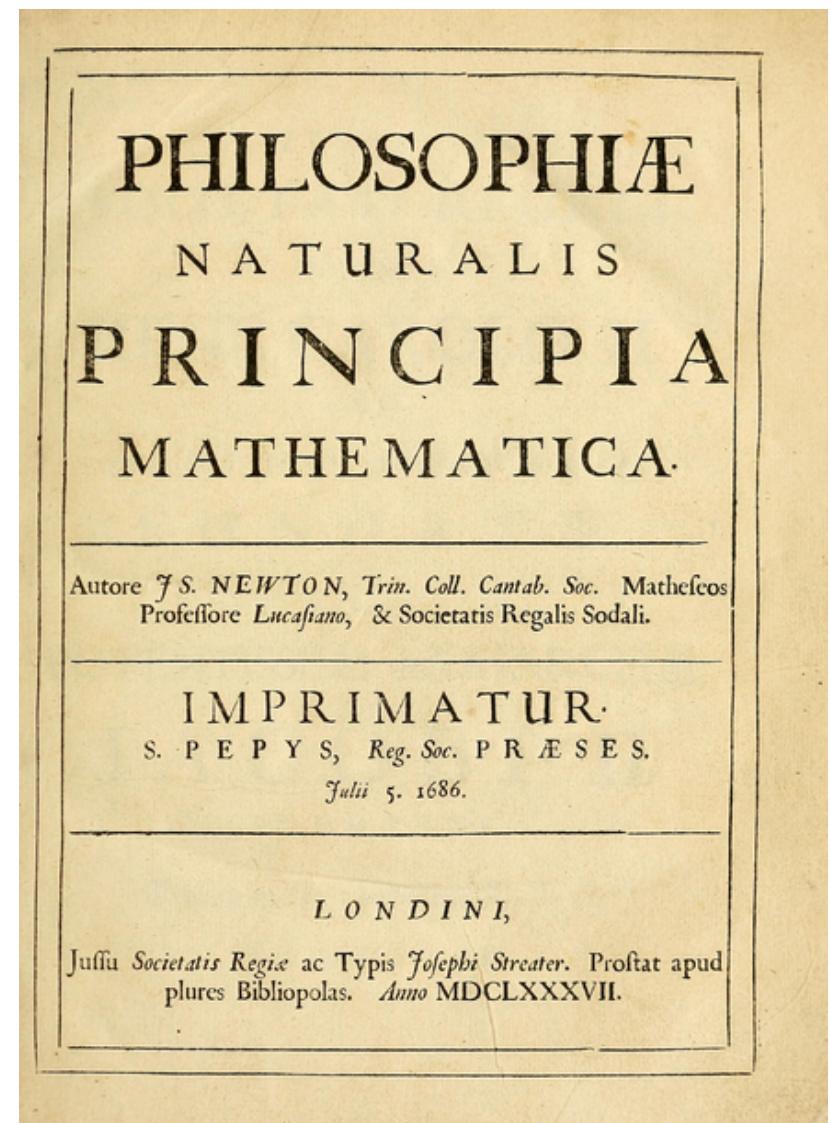


Principia Mathematica

The foundations of arithmetic in C++

Lisa Lippincott

CppCon, September 2022

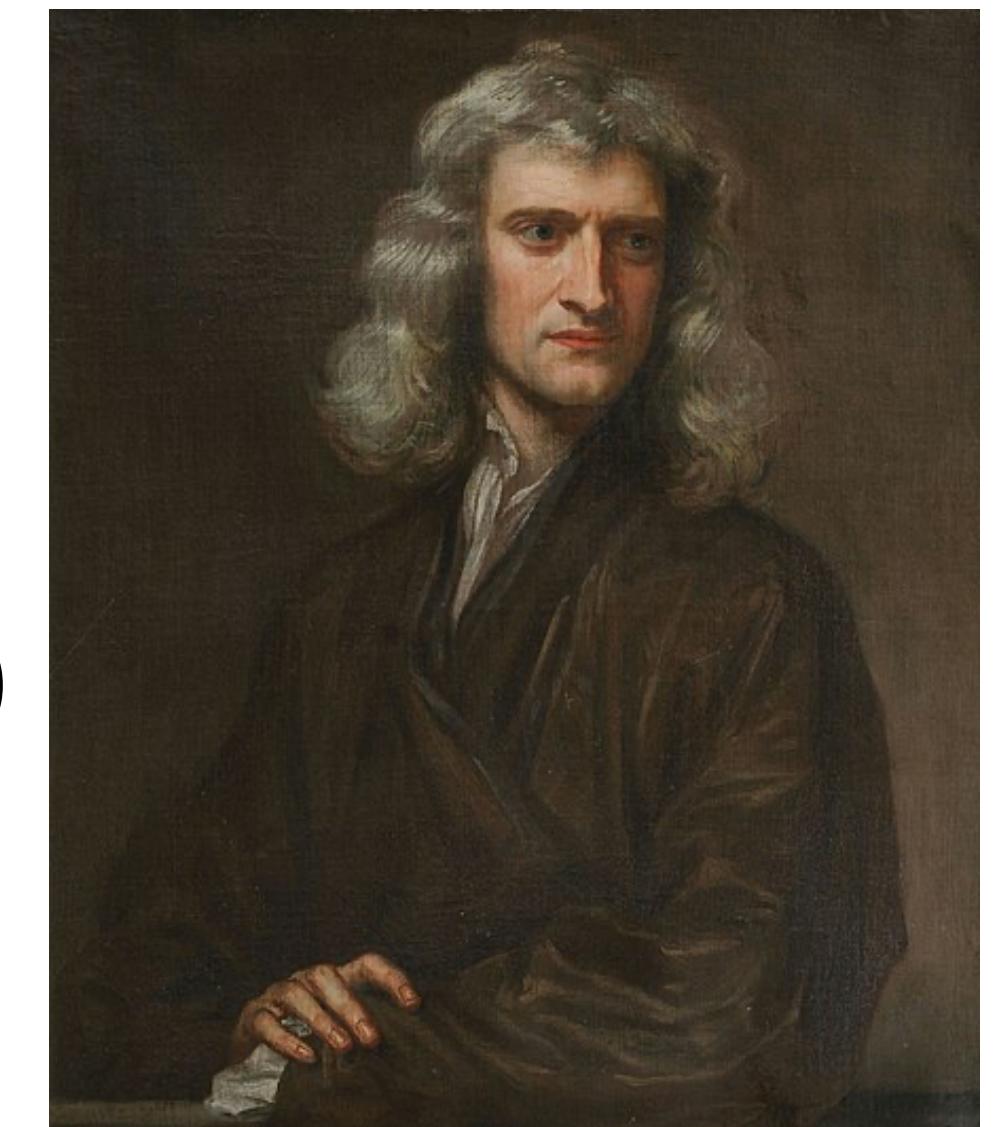


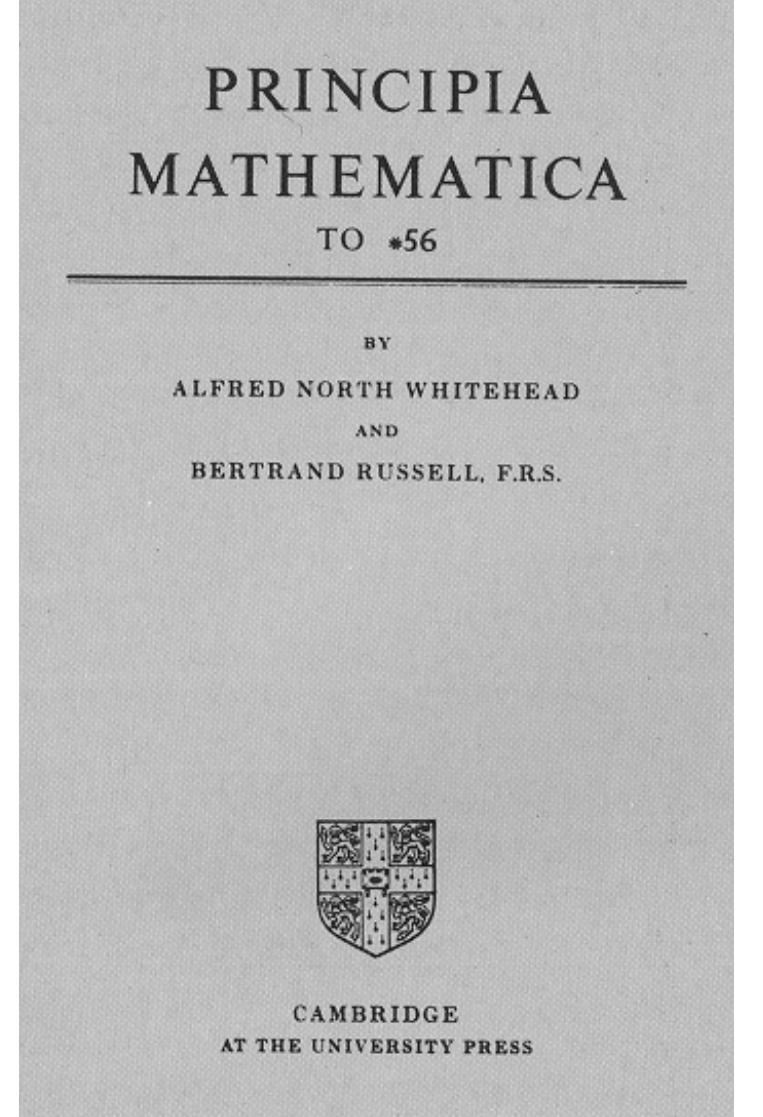
Philosophiae Naturalis Principia Mathematica

(Mathematical Principles of Natural Philosophy)

1687

Isaac Newton
1643-1727 (Gregorian)



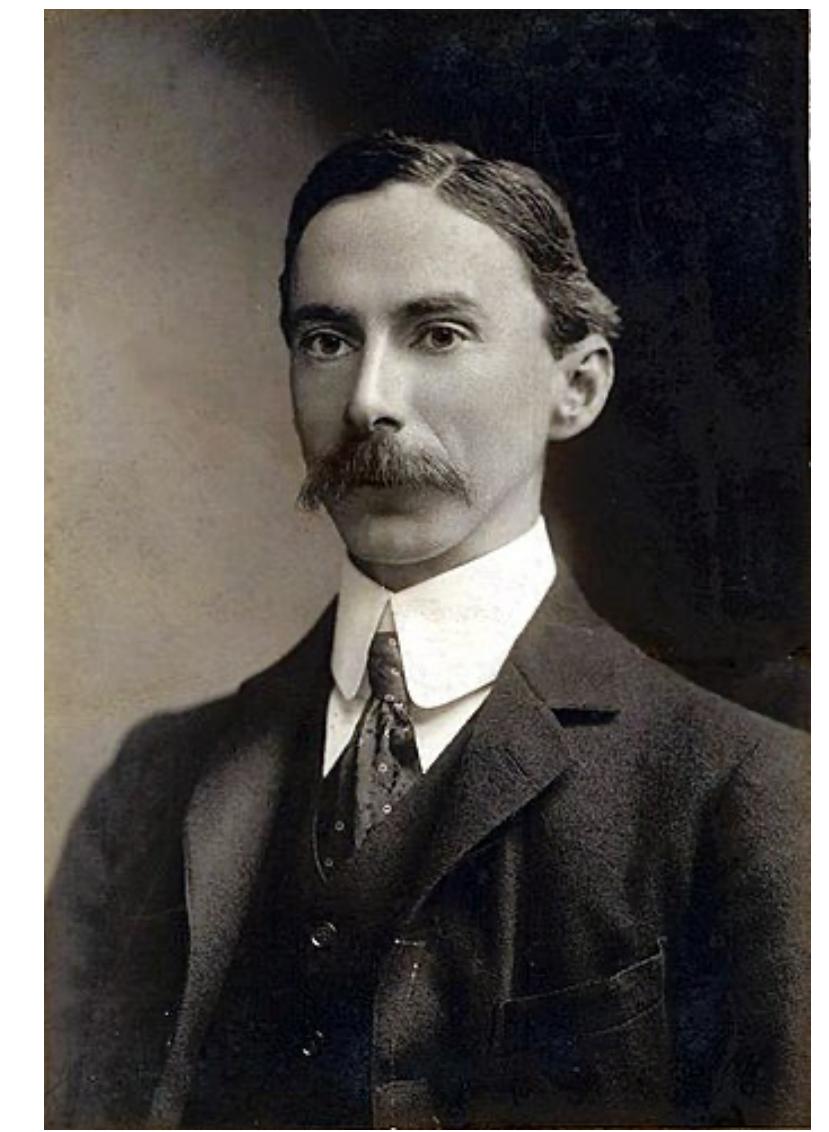


Alfred North Whitehead
1861-1947



Principia Mathematica
(Principles of Mathematics)
1910, 1912, 1913

Bertrand Russell
1872-1970



*110·632. $\vdash : \mu \in NC . \mathcal{D} . \mu +_e 1 = \hat{\xi} \{ (\exists y) . y \in \xi . \xi - t^y \in sm^{\mu} \mu \}$

Dem.

$\vdash . *110\cdot631 . *51\cdot211\cdot22 . \mathcal{D}$

$\vdash : Hp . \mathcal{D} . \mu +_e 1 = \hat{\xi} \{ (\exists y, \gamma) . y \in sm^{\mu} \mu . y \in \xi . \gamma = \xi - t^y \}$

[*13·195] $= \hat{\xi} \{ (\exists y) . y \in \xi . \xi - t^y \in sm^{\mu} \mu \} : \mathcal{D} \vdash . Prop$

*110·64. $\vdash . 0 +_e 0 = 0$ [*110·62]

*110·641. $\vdash . 1 +_e 0 = 0 +_e 1 = 1$ [*110·51·61 . *101·2]

*110·642. $\vdash . 2 +_e 0 = 0 +_e 2 = 2$ [*110·51·61 . *101·31]

*110·643. $\vdash . 1 +_e 1 = 2$

Dem.

$\vdash . *110\cdot632 . *101\cdot21\cdot28 . \mathcal{D}$

$\vdash . 1 +_e 1 = \hat{\xi} \{ (\exists y) . y \in \xi . \xi - t^y \in 1 \}$

[*54·3] $= 2 . \mathcal{D} \vdash . Prop$

The above proposition is occasionally useful. It is used at least three times, in *113·66 and *120·123·472.

*110·7·71 are required for proving *110·72, and *110·72 is used in *117·3, which is a fundamental proposition in the theory of greater and less.

*110·7. $\vdash : \beta \subset \alpha . \mathcal{D} . (\exists \mu) . \mu \in NC . Ne^{\mu} \alpha = Ne^{\mu} \beta +_e \mu$

Dem.

$\vdash . *24\cdot411\cdot21 . \mathcal{D} \vdash : Hp . \mathcal{D} . \alpha = \beta \cup (\alpha - \beta) . \beta \cap (\alpha - \beta) = \Lambda .$

[*110·32] $\mathcal{D} . Ne^{\mu} \alpha = Ne^{\mu} \beta +_e Ne^{\mu} (\alpha - \beta) : \mathcal{D} \vdash . Prop$

*110·71. $\vdash : (\exists \mu) . Ne^{\mu} \alpha = Ne^{\mu} \beta +_e \mu . \mathcal{D} . (\exists \delta) . \delta sm \beta . \delta \subset \alpha$

Dem.

$\vdash . *100\cdot3 . *110\cdot4 . \mathcal{D}$

$\vdash : Ne^{\mu} \alpha = Ne^{\mu} \beta +_e \mu . \mathcal{D} . \mu \in NC - t^{\mu} \Lambda$ (1)

$\vdash . *110\cdot3 . \mathcal{D} \vdash : Ne^{\mu} \alpha = Ne^{\mu} \beta +_e Ne^{\mu} \gamma . \equiv . Ne^{\mu} \alpha = Ne^{\mu} (\beta + \gamma) .$

[*100·3·31] $\mathcal{D} . \alpha sm (\beta + \gamma) .$

[*73·1] $\mathcal{D} . (\exists R) . R \in 1 \rightarrow 1 . D^R = \alpha . D^R = \downarrow \Lambda_{\gamma} "t^{\gamma} \beta \cup \Lambda_{\beta} \downarrow "t^{\gamma} \gamma .$

[*37·15] $\mathcal{D} . (\exists R) . R \in 1 \rightarrow 1 . \downarrow \Lambda_{\gamma} "t^{\gamma} \beta \subset D^R . R \downarrow \Lambda_{\gamma} "t^{\gamma} \beta \subset \alpha .$

[*110·12·*73·22] $\mathcal{D} . (\exists \delta) . \delta \subset \alpha . \delta sm \beta$ (2)

$\vdash . (1) . (2) . \mathcal{D} \vdash . Prop$

Volume II, page 86
(1st edition)

$$*110\cdot641. \vdash. 1 +_e 0 = 0 +_e 1 = 1 \quad [*110\cdot51\cdot61, *101\cdot2]$$

Volume II, page 86
(1st edition)

$$*110\cdot642. \vdash. 2 +_e 0 = 0 +_e 2 = 2 \quad [*110\cdot51\cdot61, *101\cdot31]$$

$$*110\cdot643. \vdash. 1 +_e 1 = 2$$

Dem.

$$\vdash. *110\cdot632. *101\cdot21\cdot28. \mathbb{D}$$

$$\vdash. 1 +_e 1 = \hat{\xi}\{(\exists y). y \in \xi. \xi - \iota^y y \in 1\}$$

$$[*54\cdot3] = 2. \mathbb{D} \vdash. \text{Prop}$$

The above proposition is occasionally useful. It is used at least three times, in *113·66 and *120·123·472.

*110·7·71 are required for proving *110·72, and *110·72 is used in *117·3, which is a fundamental proposition in the theory of greater and less.

$$*110\cdot7. \vdash: \beta \subset \alpha. \mathbb{D}. (\exists \mu). \mu \in NC. Nc^e \alpha = Nc^e \beta +_e \mu$$

The foundations of arithmetic in C++

The foundations of (arithmetic in C++)

(The foundations of arithmetic) in C++

result_type function_name(parameter_list)

interface

{

// preconditions...

The calling function is responsible
for the top part of the interface.

implementation;

// postconditions...

The called function is responsible
for the bottom part of the interface.

}

+ b	a + b	a & b	a == b
- b	a - b	a b	a < b
	a * b	a ^ b	a > b
~b	a / b		a <= b
	a % b	a << b	a >= b
		a >> b	a != b
			a <=> b

++ b	a += b	a &= b	a = b
-- b	a -= b	a = b	
	a *= b	a ^= b	
a ++	a /= b		
a --	a %= b	a <<= b	
		a >>= b	

Stability

Over certain periods, an object's state, and therefore its value, remains stable.

Substitutability

At certain times, two different objects of the same type have interchangeable values.

Repeatability

An operation may be repeated by a sufficiently similar operation, producing similar results.

$+ b$	$a + b$	$a \& b$	$a == b$
$- b$	$a - b$	$a b$	$a < b$
	$a * b$	$a ^ b$	$a > b$
$\sim b$	a / b		$a \leq b$
	$a \% b$	$a \ll b$	$a \geq b$
		$a \gg b$	$a != b$
			$a \Leftrightarrow b$

$++ b$	$a += b$	$a \&= b$	$a = b$
$-- b$	$a -= b$	$a = b$	
	$a *= b$	$a ^= b$	
$a ++$	$a /= b$		
$a --$	$a \%= b$	$a \ll= b$	
		$a \gg= b$	

Right of stability

A right of stability is transferred from the caller to the implementation on entry, and from the implementation to the caller on exit.

+ b	a + b	a & b	a == b
- b	a - b	a b	a < b
~ b	a * b	a ^ b	a > b
	a / b	a << b	a <= b
	a % b	a >> b	a >= b
			a != b
			a <=> b

Immunity from instability

The caller extends immunity from instability to the implementation for the duration of the operation.

++ b	a += b	a &= b	a = b
-- b	a -= b	a = b	
	a *= b	a ^= b	
	a /= b	a <= b	
	a %= b	a >= b	
		a <<= b	
		a >>= b	

Right of stability

A right of stability is transferred from the caller to the implementation on entry, and from the implementation to the caller on exit.

+ b

- b

\sim b

a + b
a - b
a * b

a / b
a % b

a & b
a | b
a ^ b

a << b
a >> b

a == b
a < b
a > b

a <= b
a >= b
a != b
a <=> b

++ b

-- b

a ++
a --

a += b

a -= b

a *= b

a /= b

a %= b

a &= b

a |= b

a ^= b

a <<= b

a >>= b

a = b

Right of stability

A right of stability of the function result is transferred from the implementation to the caller.

+ b	a + b	a & b	a == b
- b	a - b	a b	a < b
	a * b	a ^ b	a > b
~b	a / b		a <= b
	a % b	a << b	a >= b
		a >> b	a != b
			a <=> b

Substitutability

Most operations here that modify an argument return a reference aliased to that argument.

++ b	a += b
-- b	a -= b
	a *= b
	a /= b
a ++	a %= b
a --	

a &= b	a = b
a = b	
a ^= b	
	a <= b
	a >= b

claim substitutable(&result, &a);

claim substitutable(&result, &b);

$+ b$	$a + b$	$a \& b$	<code>a == b</code>	if (result) claim substitutable(a, b);
$- b$	$a - b$	$a b$	<code>a < b</code>	
			<code>a > b</code>	
$\sim b$	a / b		<code>a <= b</code>	
			<code>a >= b</code>	
	$a \% b$	$a << b$	<code>a >> b</code>	
			<code>a != b</code>	if (!result) claim substitutable(a, b);
			<code>a <=> b</code>	

$++ b$	$a += b$	$a \&= b$	<code>a = b</code>	claim substitutable(a, b);
$-- b$	$a -= b$	$a = b$		
		$a *= b$	$a ^= b$	
$a ++$	$a /= b$			
$a --$	$a \%= b$	$a <= b$		
		$a >= b$		

Substitutability

These operations may also have substitutability as a postcondition.

+ b	a + b	a & b	a == b
- b	a - b	a b	a < b
~ b	a * b	a ^ b	a > b
	a / b		a <= b
	a % b	a << b	a >= b
		a >> b	a != b
			a <=> b

Discernible input

Most parameters are part of the discernible input to the operations.

An operation is repeated when these parameters are repeated.

++ b	a += b	a &= b	a = b
-- b	a -= b	a = b	
	a *= b	a ^= b	
	a /= b		
	a %= b	a <= b	
		a >= b	

The left parameter of operator= is not part of the discernible input.

+ b	a + b	a & b	a == b
- b	a - b	a b	a < b
	a * b	a ^ b	a > b
~b	a / b		a <= b
	a % b	a << b	a >= b
		a >> b	a != b
			a <=> b

Discernible output

The function result is part of the discernible output for every operation listed here.

Left-hand arguments of `operator++` and `operator--` are also part of the discernible output.

++ b	a += b	a &= b	a = b
-- b	a -= b	a = b	
	a *= b	a ^= b	
a++	a /= b		
a--	a %= b	a <= b	
		a >= b	

If an operation is repeated, the discernible output will be repeated.

```
int& operator=( int& a, const int b )
```

```
interface
```

```
{
```

```
claim_right a;
```

// right of stability

```
claim_immunity b;
```

// immunity from instability

```
discern b;
```

// discernible input

```
implementation;
```

```
claim substitutable( &a, &result );
```

// substitutability

```
claim_right result;
```

// right of stability

```
discern result;
```

// discernible output

```
claim substitutable( result, b );
```

// substitutability

```
}
```

bool

unsigned char

unsigned short int

unsigned int

unsigned long int

unsigned long long int

char

char8_t

char16_t

char32_t

wchar_t

signed char

short int

int

long int

long long int

bool

unsigned char

unsigned short int

unsigned int

unsigned long int

unsigned long long int

unsigned longer long int

unsigned extremely long int

unsigned overwhelmingly long int

unsigned mind numbingly vast int

unsigned oppressively colossal int

char

char8_t

char16_t

char32_t

wchar_t

signed char

short int

int

long int

long long int

longer long int

extremely long int

overwhelmingly long int

mind numbingly vast int

oppressively colossal int

Unsigned

$0 \leq \text{value} < 2^{\text{width}}$

unsigned oppressively colossal int

← width →

Signed

$-2^{\text{width}-1} \leq \text{value} < 2^{\text{width}-1}$

oppressively colossal int

← width →

Unsigned

$$0 \leq \text{value} < 2^{\text{width}}$$

unsigned oppressively colossal int



Signed

$$-2^{\text{width}-1} \leq \text{value} < 2^{\text{width}-1}$$

oppressively colossal int



```
class integer_kind
{
    // ...
    constexpr bool      is_signed() const;
    constexpr bit_size_t width()      const;
};
```

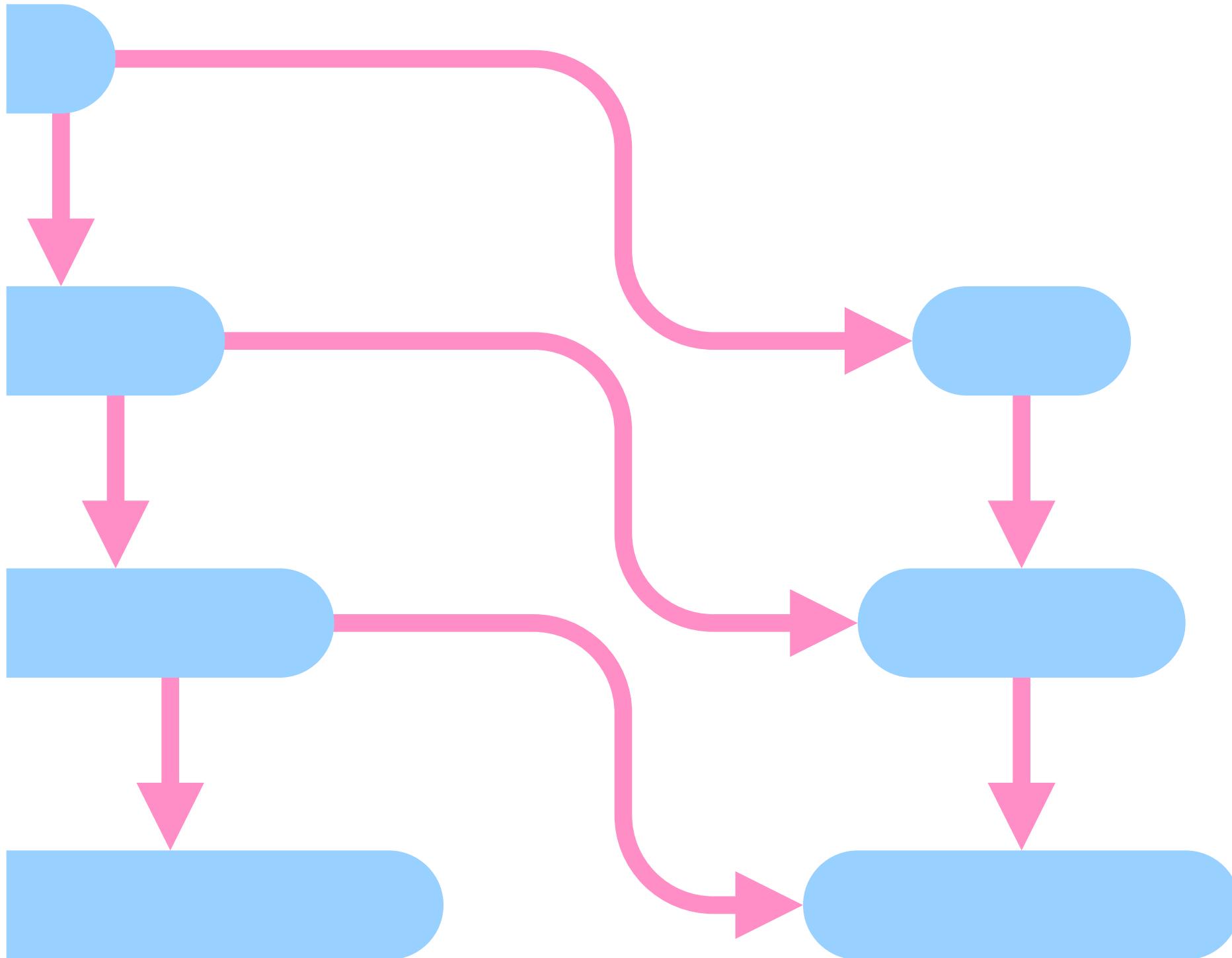
```
class integer_kind
{
// ...
constexpr bool      is_signed() const;
constexpr bit_size_t width()      const;
};
```

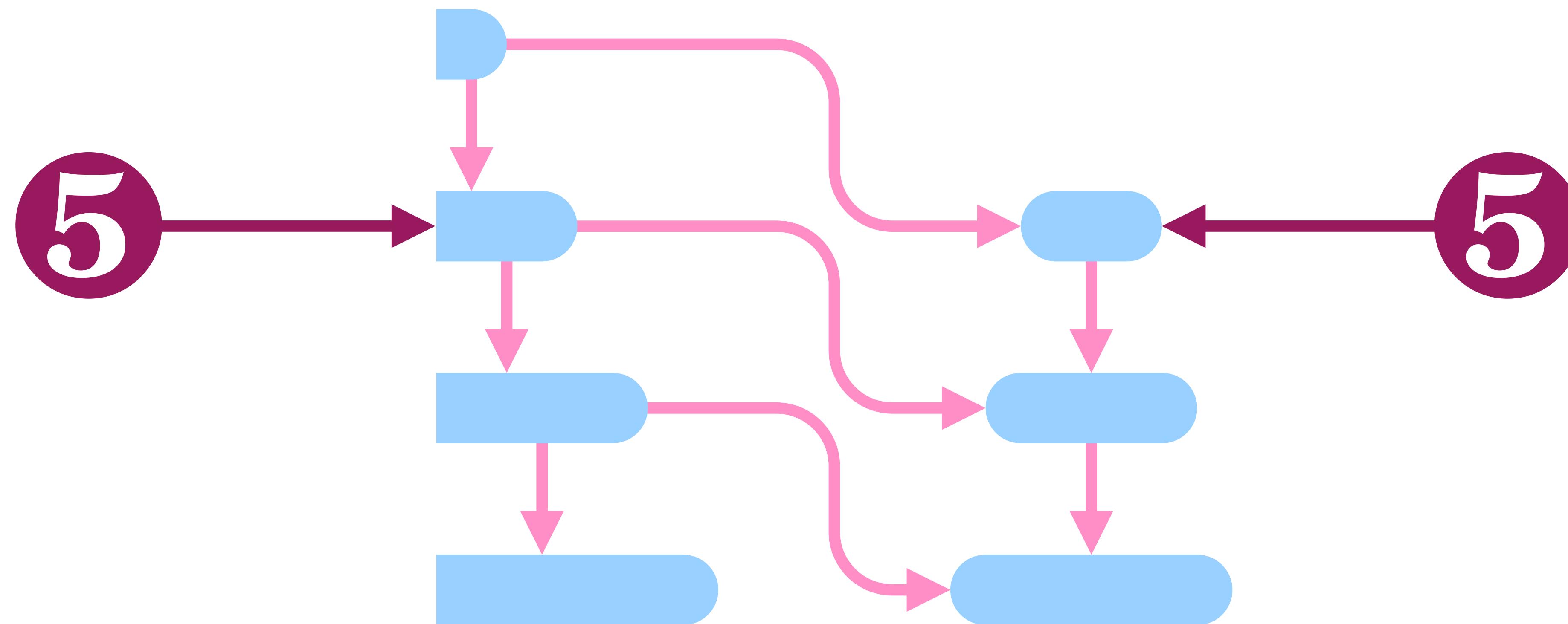
```
inline constexpr
bool operator<( integer_kind a, integer_kind b )
{
    return a.is_signed() <= b.is_signed()
        && a.width()      <  b.width();
}
```

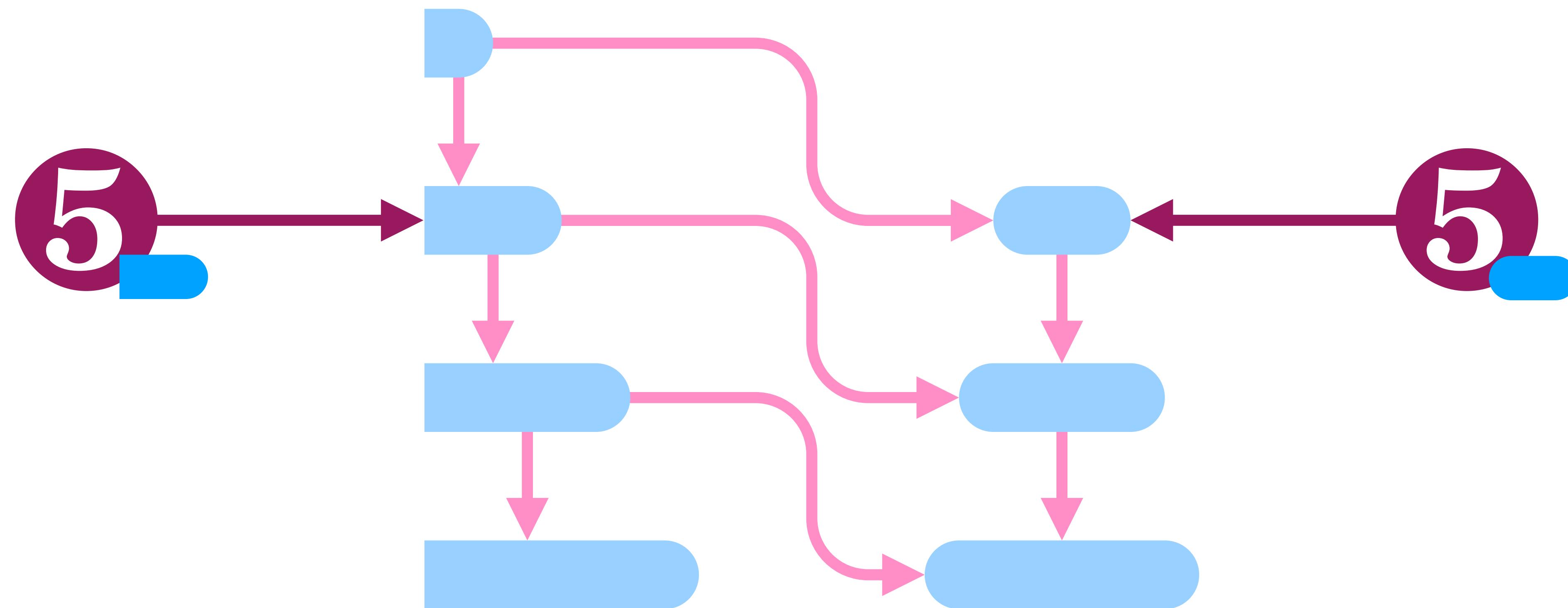
```
class integer_kind
{
    // ...
    constexpr bool      is_signed() const;
    constexpr bit_size_t width()     const;
};
```

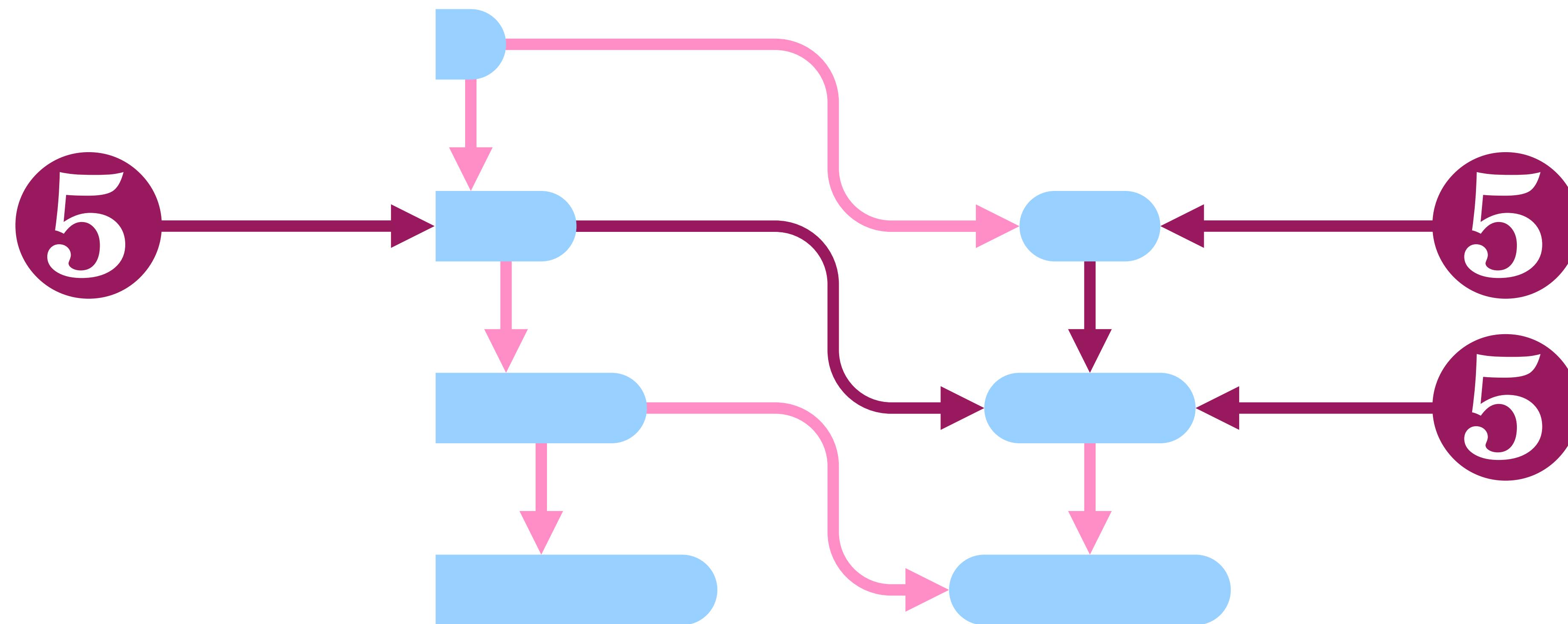
```
inline constexpr
bool operator<( integer_kind a, integer_kind b )
{
    return a.is_signed() <= b.is_signed()
        && a.width()       <   b.width();
}
```

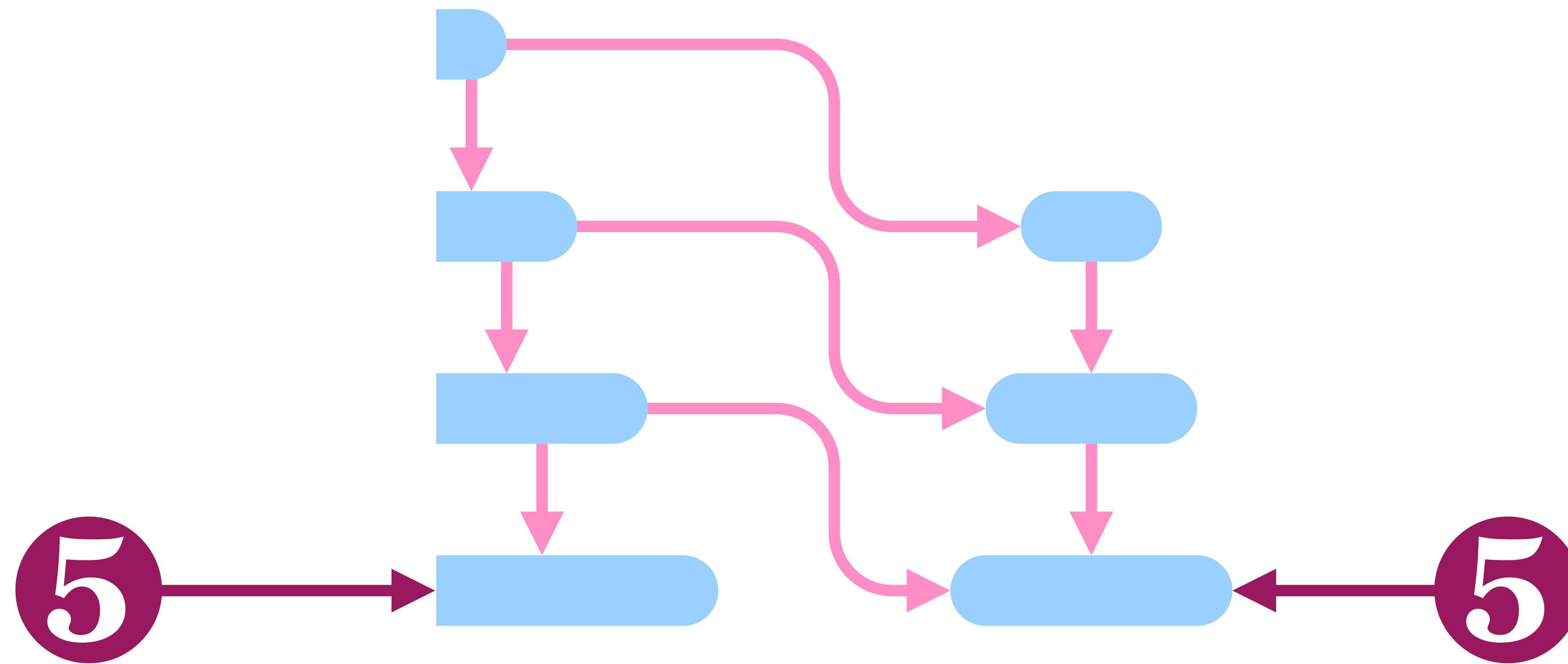
```
template < class To, class From >
    requires ( integer_kind_of<To> >= integer_kind_of<From> )
    To convert( const From& );
```

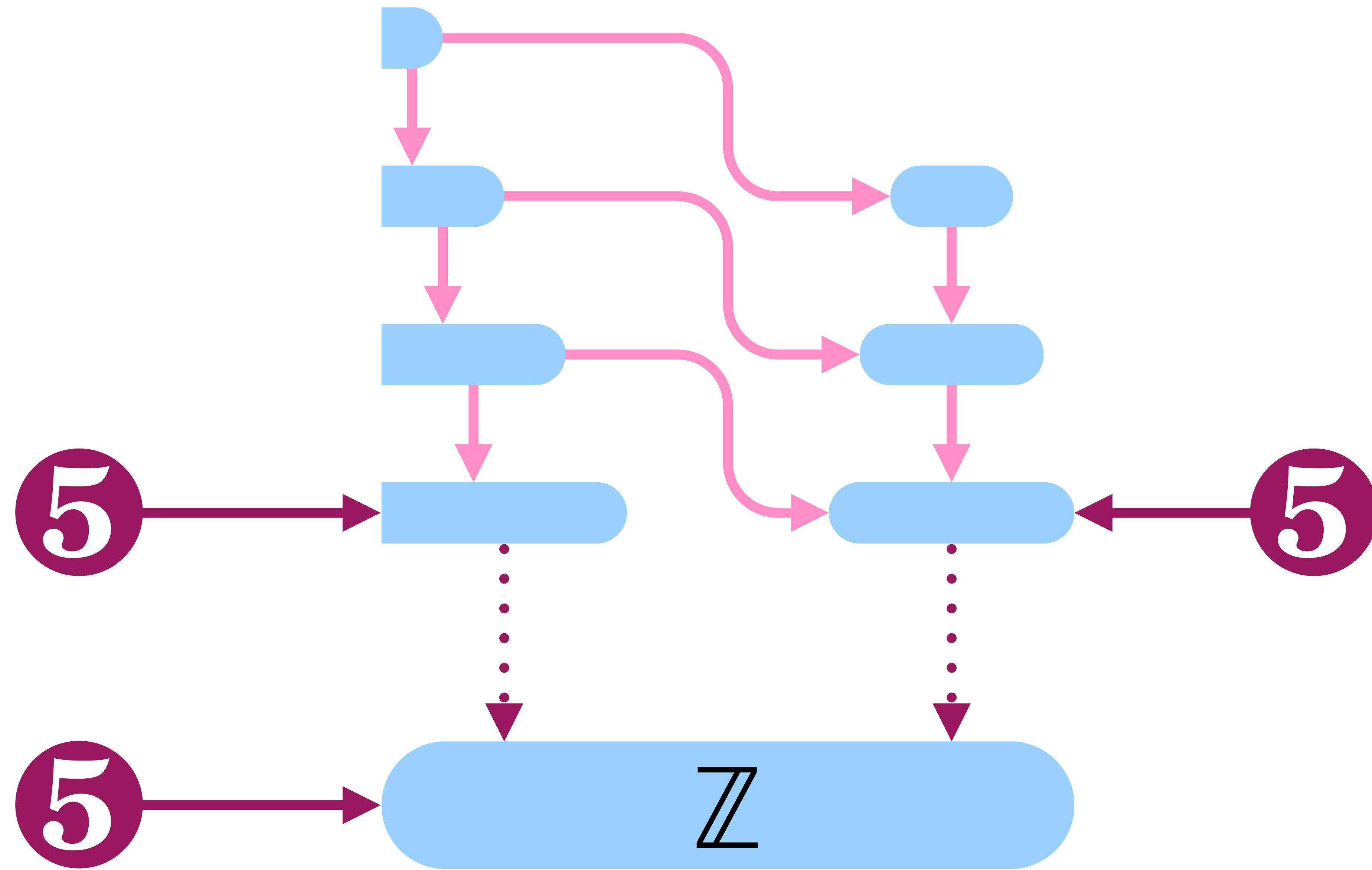












- To convert(const From&); ————— Strong type requirement:
All values must be convertible
- To convert_narrowing(const From&); ————— Strong precondition:
Argument must be convertible
- To convert_modular(const From&); ————— Weak postcondition:
Only the low bits are converted

```
int& operator=( int& a, const int b )
```

```
interface
```

```
{
```

```
claim_right a;
```

// right of stability

```
claim_immunity b;
```

// immunity from instability

```
discern b;
```

// discernible input

```
implementation;
```

```
claim substitutable( &a, &result );
```

// substitutability

```
claim_right result;
```

// right of stability

```
discern result;
```

// discernible output

```
claim substitutable( result, b );
```

// substitutability

```
}
```

```
template < std::integral A, std::integral B >
```

```
A& operator=( A& a, const B b )
```

```
interface
```

```
{
```

```
claim_right a;
```

// right of stability

```
claim_immunity b;
```

// immunity from instability

```
discern b;
```

// discernible input

```
implementation;
```

```
claim substitutable( &a, &result );
```

// substitutability

```
claim_right result;
```

// right of stability

```
discern result;
```

// discernible output

```
claim substitutable( result, convert_modular<A>( b ) );
```

```
}
```

+ b	a + b	a & b	a == b	convert
- b	a - b	a b	a < b	convert_modular
	a * b	a ^ b	a > b	convert_narrowing
~b	a / b		a <= b	
	a % b	a << b	a >= b	
		a >> b	a != b	
			a<=>b	

++ b	a += b	a &= b	a = b
-- b	a -= b	a = b	
	a *= b	a ^= b	
a++	a /= b		
a--	a %= b	a<<=b	
		a>>=b	

<code>+ b</code>	<code>a + b</code>	<code>a & b</code>	<code>a == b</code>	<code>convert</code>
<code>- b</code>	<code>a - b</code>	<code>a b</code>	<code>a < b</code>	<code>convert_modular</code>
	<code>a * b</code>	<code>a ^ b</code>	<code>a > b</code>	<code>convert_narrowing</code>
<code>~b</code>	<code>a / b</code>		<code>a <= b</code>	
	<code>a % b</code>	<code>a << b</code>	<code>a >= b</code>	
		<code>a >> b</code>	<code>a != b</code>	
			<code>a <=> b</code>	

<code>++ b</code>	<code>a += b</code>	<code>a &= b</code>
<code>-- b</code>	<code>a -= b</code>	<code>a = b</code>
	<code>a *= b</code>	<code>a ^= b</code>
<code>a++</code>	<code>a /= b</code>	
<code>a--</code>	<code>a %= b</code>	<code>a <= b</code>
		<code>a >= b</code>

The behavior of an expression of the form `E1 op= E2` is equivalent to `E1 = E1 op E2` except that `E1` is evaluated only once.

7.6.19 [expr.ass]

```
template < std::integral A, std::integral B >
A& operator+=( A& a, const B b )
interface
```

```
{  
    claim_right a;                                // right of stability  
    claim_immunity b;                            // immunity from instability  
    discern a;                                    // discernible input  
    discern b;                                    // discernible input
```

```
const auto expected_result = convert_modular<A>( a+b );
```

implementation:

```
claim substitutable( &a, &result );                // substitutability  
claim_right result;                            // right of stability  
discern result;                                // discernible output
```

```
claim substitutable( result, expected_result );    // substitutability  
}
```

```
template < std::integral A, std::integral B >
A& operator+=( A& a, const B b )
interface
```

```
{  
    claim_right a;                                // right of stability  
    claim_immunity b;                            // immunity from instability  
    discern a;                                    // discernible input  
    discern b;                                    // discernible input
```

```
const auto expected_result = convert_modular<A>( a+b );
```

implementation:

```
claim substitutable( &a, &result );                // substitutability  
claim_right result;                            // right of stability  
discern result;                                // discernible output
```

```
claim result == expected_result;                // substitutability  
}
```

$+ b$	$a + b$	$a \& b$	$a == b$	convert
$- b$	$a - b$	$a b$	$a < b$	convert_modular
	$a * b$	$a ^ b$	$a > b$	convert_narrowing
$\sim b$	a / b		$a \leq b$	
	$a \% b$	$a \ll b$	$a \geq b$	
		$a \gg b$	$a != b$	
			$a \Leftrightarrow b$	

$++ b$	$a += b$	$a \&= b$
$-- b$	$a -= b$	$a = b$
	$a *= b$	$a ^= b$
$a ++$	$a /= b$	
$a --$	$a \% = b$	$a \ll= b$
		$a \gg= b$

The behavior of an expression of the form $E1 \ op= E2$ is equivalent to $E1 = E1 \ op E2$ except that $E1$ is evaluated only once.

7.6.19 [expr.ass]

+ b	a + b	a & b	a == b	convert
- b	a - b	a b	a < b	convert_modular
	a * b	a ^ b	a > b	convert_narrowing
~b	a / b		a <= b	
	a % b	a << b	a >= b	
next		a >> b	a != b	
			a<=>b	

++ b

const auto expected_result = next(b);

-- b

a ++

a --

+ b	a + b	a & b	a == b	convert
- b	a - b	a b	a < b	convert_modular
	a * b	a ^ b	a > b	convert_narrowing
~b	a / b		a <= b	
	a % b	a << b	a >= b	
next		a >> b	a != b	
prev			a<=>b	

-- b

const auto expected_result = prev(b);

a++

a--

+ b	a + b	a & b	a == b	convert
- b	a - b	a b	a < b	convert_modular
	a * b	a ^ b	a > b	convert_narrowing
~b	a / b		a <= b	
	a % b	a << b	a >= b	
next		a >> b	a != b	
prev			a<=>b	

```
const auto expected_result = a;
const auto expected_a      = next( a );
```

a++

// ...implementation...

a--

```
claim result == expected_result;
claim a      == expected_a;
```

+ b	a + b	a & b	a == b	convert
- b	a - b	a b	a < b	convert_modular
	a * b	a ^ b	a > b	convert_narrowing
~b	a / b		a <= b	
	a % b	a << b	a >= b	
next		a >> b	a != b	
prev			a<=>b	

```
const auto expected_result = a;
const auto expected_a      = prev( a );
```

// ...implementation...

a --

```
claim result == expected_result;
claim a      == expected_a;
```

bool

unsigned char

unsigned short int

unsigned int

unsigned long int

unsigned long long int

unsigned longer long int

unsigned extremely long int

unsigned overwhelmingly long int

unsigned mind numbingly vast int

unsigned oppressively colossal int

char

char8_t

char16_t

char32_t

wchar_t

signed char

short int

int

long int

long long int

longer long int

extremely long int

overwhelmingly long int

mind numbingly vast int

oppressively colossal int

bool

char

unsigned char

char8_t

signed char

unsigned short int

char16_t

short int

unsigned int

char32_t

int

unsigned long int

wchar_t

long int

unsigned long long int

long long int

unsigned longer long int

longer long int

unsigned extremely long int

extremely long int

unsigned overwhelmingly long int

overwhelmingly long int

unsigned mind numbingly vast int

mind numbingly vast int

unsigned oppressively colossal int

oppressively colossal int

If T is on this list:

`wchar_t`
`char16_t`
`char32_t`

T promotes to the first type in this list to which it has a nice conversion:

If T is any other integral type

T promotes to the first type in this list to which it has a nice conversion:

`int`
`unsigned int`
`long`
`unsigned long`
`long long`
`unsigned long long`
the underlying type of T

`int`
`unsigned int`
 T

```
template < std::integral T >
using promoted_type =
    std::conditional_t< !appears_in_list< T, wchar_t, char16_t, char32_t >,
                        first_nicely_convertable< T,
                                              int,           unsigned int,
                                              T >,
                        first_nicely_convertable< T,
                                              int,           unsigned int,
                                              long,          unsigned long,
                                              long long,     unsigned long long,
                                              std::underlying_type_t<T>      >>;
```

⚠ This gives the wrong answer for bit fields.

⚠ `underlying_type` doesn't work for character types.

```
template < std::integral T >
using promoted_type = decltype( +std::declval<T>() );
```

```
template < class T >
concept promoted_integral = std::integral< T >
                           && std::same_as< T, promoted_type< T > >;
```

```
template < class T >
concept unpromoted_integral = std::integral< T >
                               && !std::same_as< T, promoted_type< T > >;
```

⚠ This gives the wrong answer for bit fields.

```
template < std::integral B >
promoted_type<B> operator+( const B& b )
interface
{
    claim_immunity b;                                // immunity from instability
    discern b;                                       // discernible input

    const auto expected_result = convert< promoted_type<B> >( b );

implementation;

    claim_right result;                            // right of stability
    discern result;                               // discernible output

    claim result == expected_result;               // substitutability
}
```

+ b	a + b	a & b	a == b	convert
- b	a - b	a b	a < b	convert_modular
	a * b	a ^ b	a > b	convert_narrowing
~b	a / b		a <= b	
	a % b	a << b	a >= b	
next		a >> b	a != b	
prev			a<=>b	

```
template <unpromoted_integral B>
promoted_type<B> operator-( const B& b )
interface
{
    claim_immunity b;                                // immunity from instability
    discern b;                                       // discernible input

    const auto expected_result = -( +b );

    implementation;

    claim_right result;                            // right of stability
    discern result;                                // discernible output

    claim result == expected_result;                // substitutability
}
```



Promoted parameters

	a + b	a & b	a == b
- b	a - b	a b	a < b
	a * b	a ^ b	a > b
~b	a / b		a <= b
	a % b	a << b	a >= b
next		a >> b	a != b
prev			a<=>b



Integral parameters

convert
convert_modular
convert_narrowing

```
template < promoted_integral T >
inline auto usual_arithmetic_conversions( const T& a, const T& b )
{
    return std::pair( a, b );
}
```

```
template < promoted_integral A, promoted_integral B >
inline auto usual_arithmetic_conversions( const A& a, const B& b )
{
    // ...
}
```

```
template < promoted_integral A, promoted_integral B >
inline auto usual_arithmetic_conversions( const A& a, const B& b )
{
    constexpr ak = integer_kind_of<A>;
    constexpr bk = integer_kind_of<B>;
    if      constexpr ( ak > bk )      return std::pair( a, convert<A>(b) );
    else if constexpr ( ak < bk )      return std::pair( convert<B>(a),
                                                       b );
    else if constexpr ( ak == bk )
    {
        // Can convert either way...
    }
    else // if the integer kinds are unordered
    {
        // Cannot convert either way...
    }
}
```

```
constexpr ak = integer_kind_of<A>;
constexpr bk = integer_kind_of<B>

if      constexpr ( ak > bk )    return std::pair(               a, convert<A>(b) );
else if constexpr ( ak < bk )    return std::pair( convert<B>(a),           b );
else if constexpr ( ak == bk )
{
    // Can convert either way. Break the tie with integer conversion rank.

    if constexpr ( integer_conversion_rank_is_less< A, B > )
        return std::pair( convert<B>(a), b );
    else
        return std::pair( a, convert<A>(b) );
}

else // if the integer kinds are unordered
{
    // Cannot convert either way...
}
```

```
else if constexpr ( ak < bk )    return std::pair( convert<B>(a),  
else if constexpr ( ak == bk )  
{  
    // Can convert either way. Break the tie with integer conversion rank.  
  
    if constexpr ( integer_conversion_rank_is_less<A, B > )  
        return std::pair( convert<B>(a), b );  
    else  
        return std::pair( a, convert<A>(b) );  
}  
else // if the integer kinds are unordered  
{  
    // Cannot convert either way. Try again with unsigned arguments.  
  
    const auto unsigned_a = convert_modular< std::make_unsigned_t<A> >( a );  
    const auto unsigned_b = convert_modular< std::make_unsigned_t<B> >( b );  
  
    return usual_arithmetic_conversions( unsigned_a, unsigned_b );  
}  
}
```

```
template < std::integral A, std::integral B >
using UAC_type = decltype( std::declval<A>() + std::declval<B>() );
```

```
template < std::integral A, std::integral B >
inline auto usual_arithmetic_conversions( const A& a, const B& b )
{
    using U = UAC_type<A, B>;
    return std::pair( convert_modular<U>( a ), convert_modular<U>( b ) );
}
```

```
template < std::promoted_integral A, std::promoted_integral B >
    requires ( !std::same_as<A,B> )
UAC_type<A,B> operator+( const A& a, const B& b )
interface
{
    claim_immunity a;                                // immunity from instability
    claim_immunity b;                                // immunity from instability
    discern a;                                       // discernible input
    discern b;                                       // discernible input

    const auto [ a1, b1 ] = usual_arithmetic_conversions( a, b );
    const auto expected_result = a1 + b1;

implementation;

    claim_right_result;                            // right of stability
    discern result;                               // discernible output

    claim result == expected_result;               // substitutability
}
```



Promoted parameters

- b a << b next
~ b a >> b prev



Integral parameters

convert
convert_modular
convert_narrowing



Matching, promoted parameters

a + b	a & b	a == b
a - b	a b	a < b
a * b	a ^ b	a > b
a / b		a <= b
a % b		a >= b
		a != b
		a <=> b

Regarding a/b and $a \% b$:

If the second operand of $/$ or $\%$ is zero the behavior is undefined.

...if the quotient a/b is representable in the type of the result, $(a/b) * b + a \% b$ is equal to a ; otherwise, the behavior of both a/b and $a \% b$ is undefined.

7.6.5 [expr.mul]

Regarding $a \ll b$ and $a \gg b$:

The behavior is undefined if the right operand is negative, or greater than or equal to the width of the promoted left operand.

7.6.7 [expr.shift]

a / b

```
claim b != 0;  
// ...implementation...
```

a % b

```
claim b != 0;  
const auto quotient = a / b;  
// ...implementation...  
  
claim quotient * b + result == a;
```

a << b
a >> b

```
const auto b1 = convert_narrowing<bit_size_t>( b );  
claim b1 < integer_kind_of<A>.width();  
// ...implementation...
```

If during the evaluation of an expression, the result is not mathematically defined or not in the range of representable values for its type, the behavior is undefined.

7.1 [expr.pre]

Arithmetic for the unsigned type is performed modulo 2^N . [Note: Unsigned arithmetic does not overflow. Overflow for signed arithmetic yields undefined behavior (7.1).
—end note]

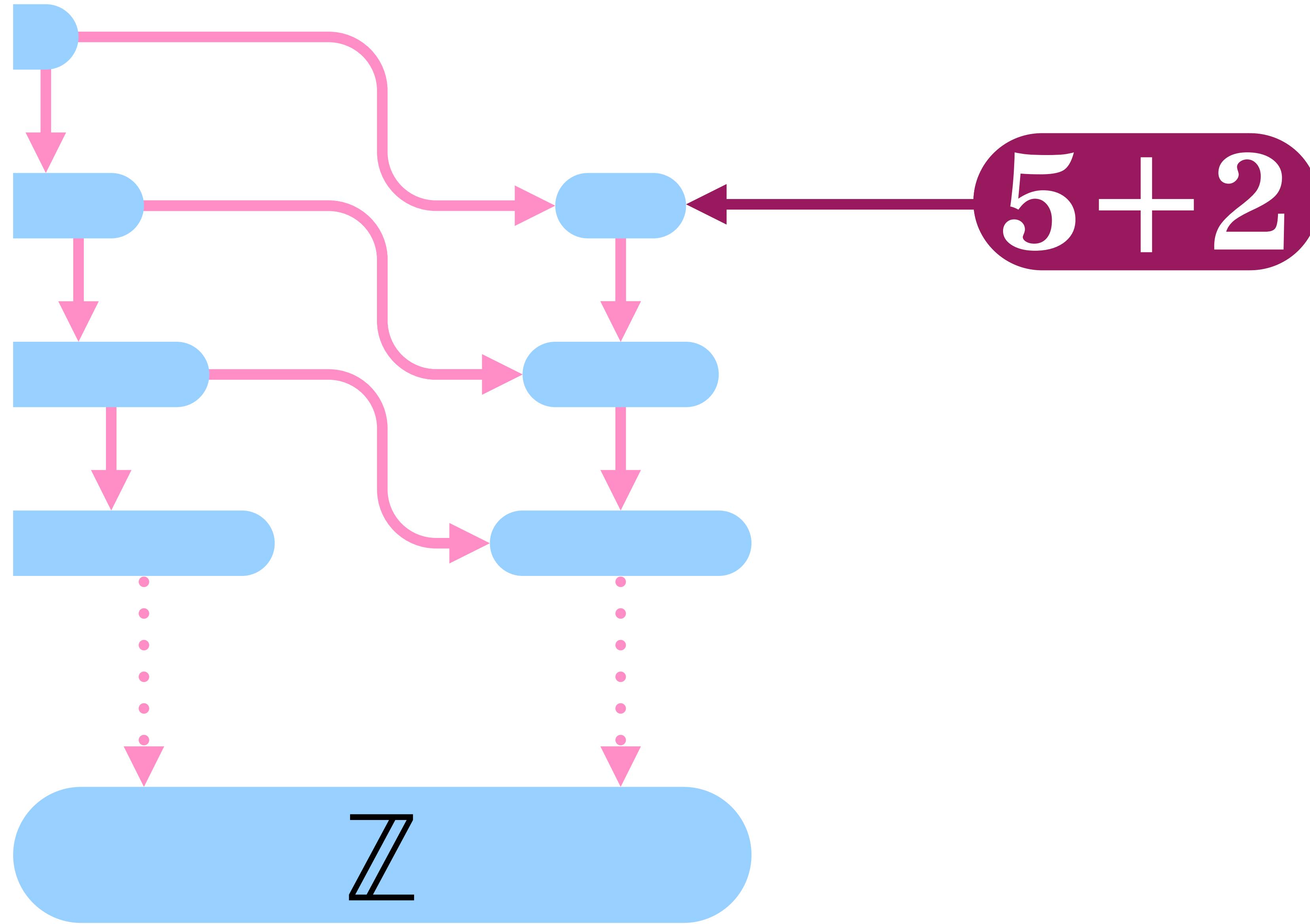
6.7.1 [basic.fundamental]

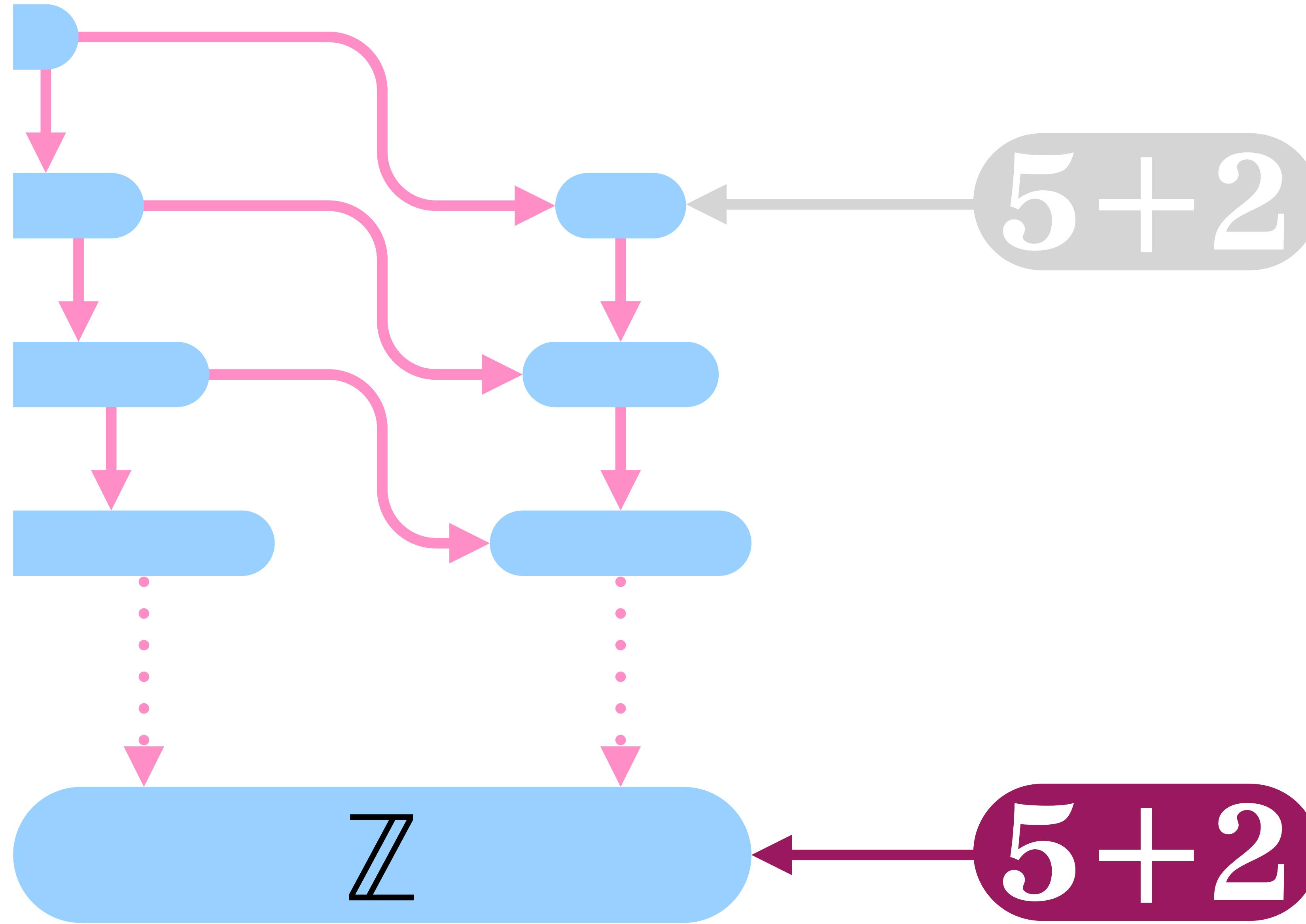
If during the evaluation of an expression, the result is not mathematically defined or not in the range of representable values for its type, the behavior is undefined.

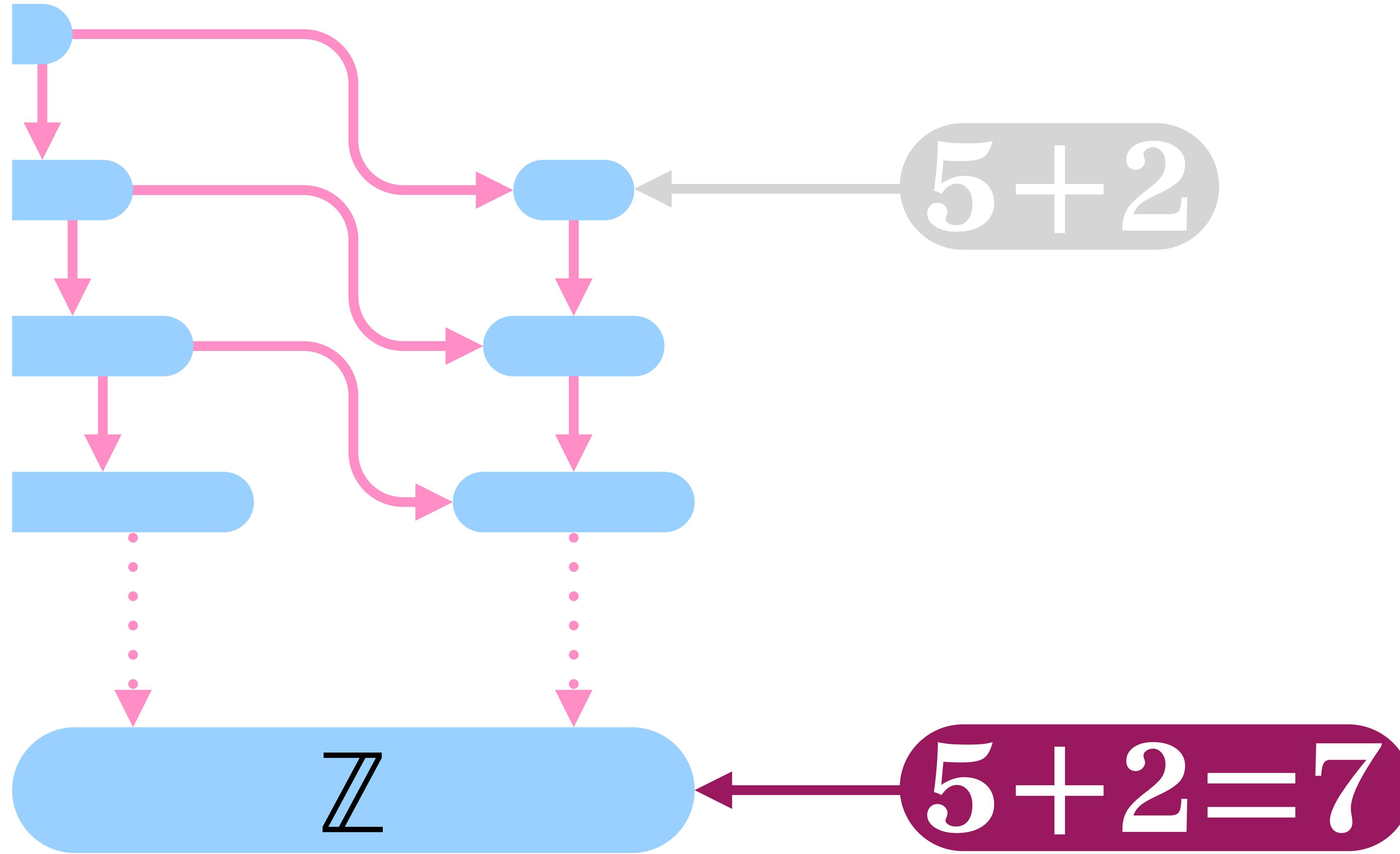
7.1 [expr.pre]

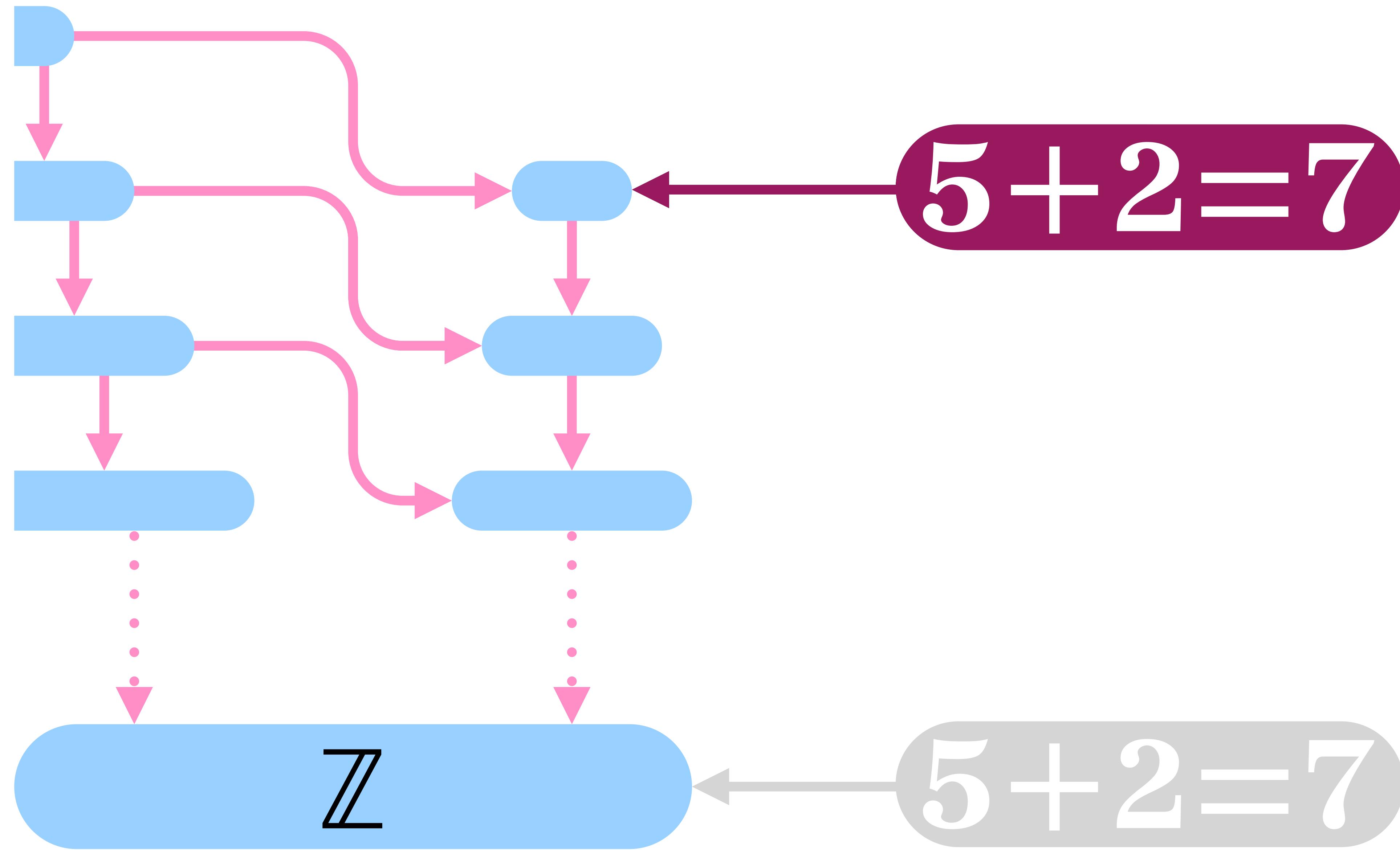
The value of **E1** \ll **E2** is the unique value congruent to $\mathbf{E1} \times 2^{\mathbf{E2}}$ modulo 2^N , where N is the width of the type of the result.

7.6.7 [expr.shift]

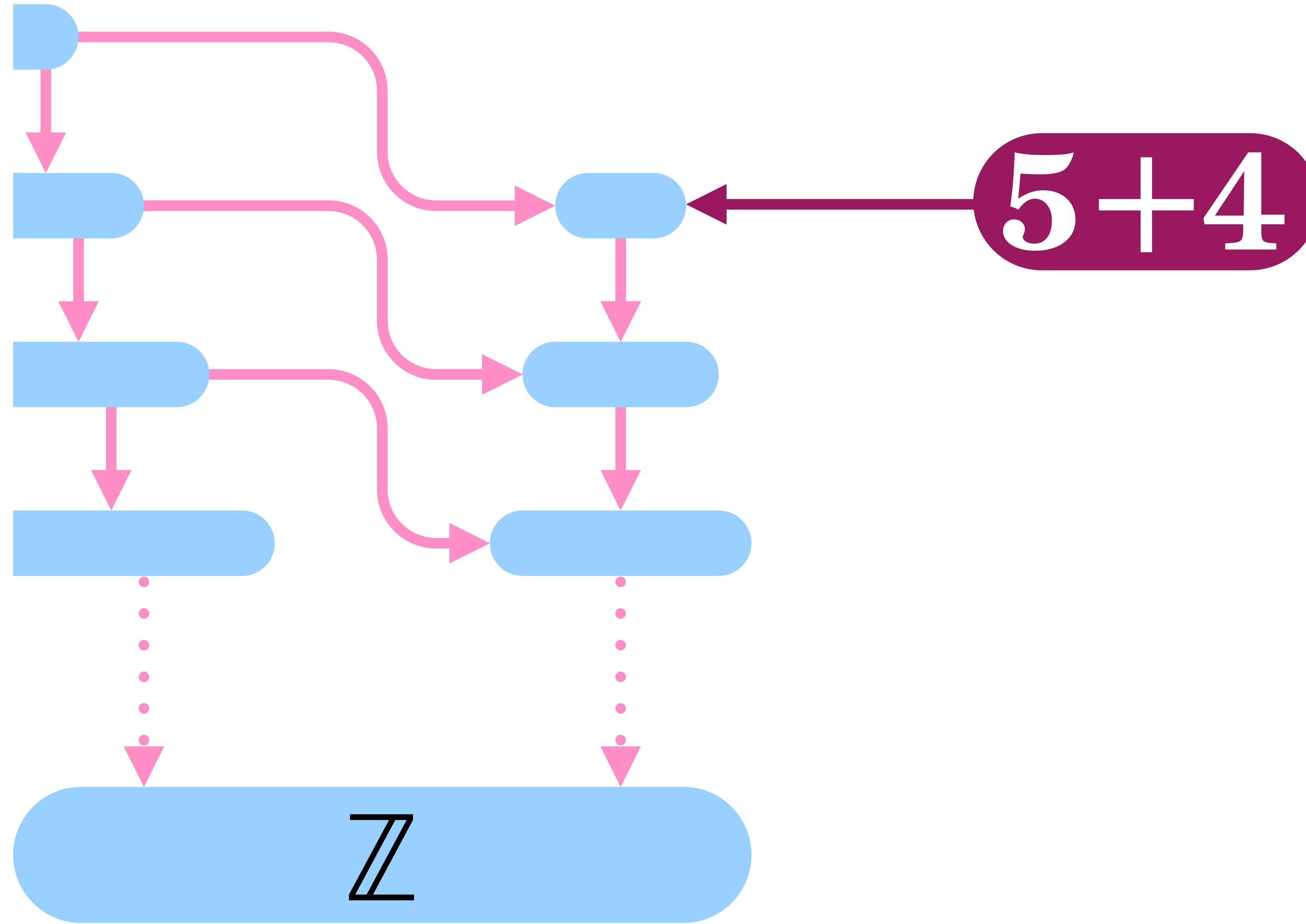


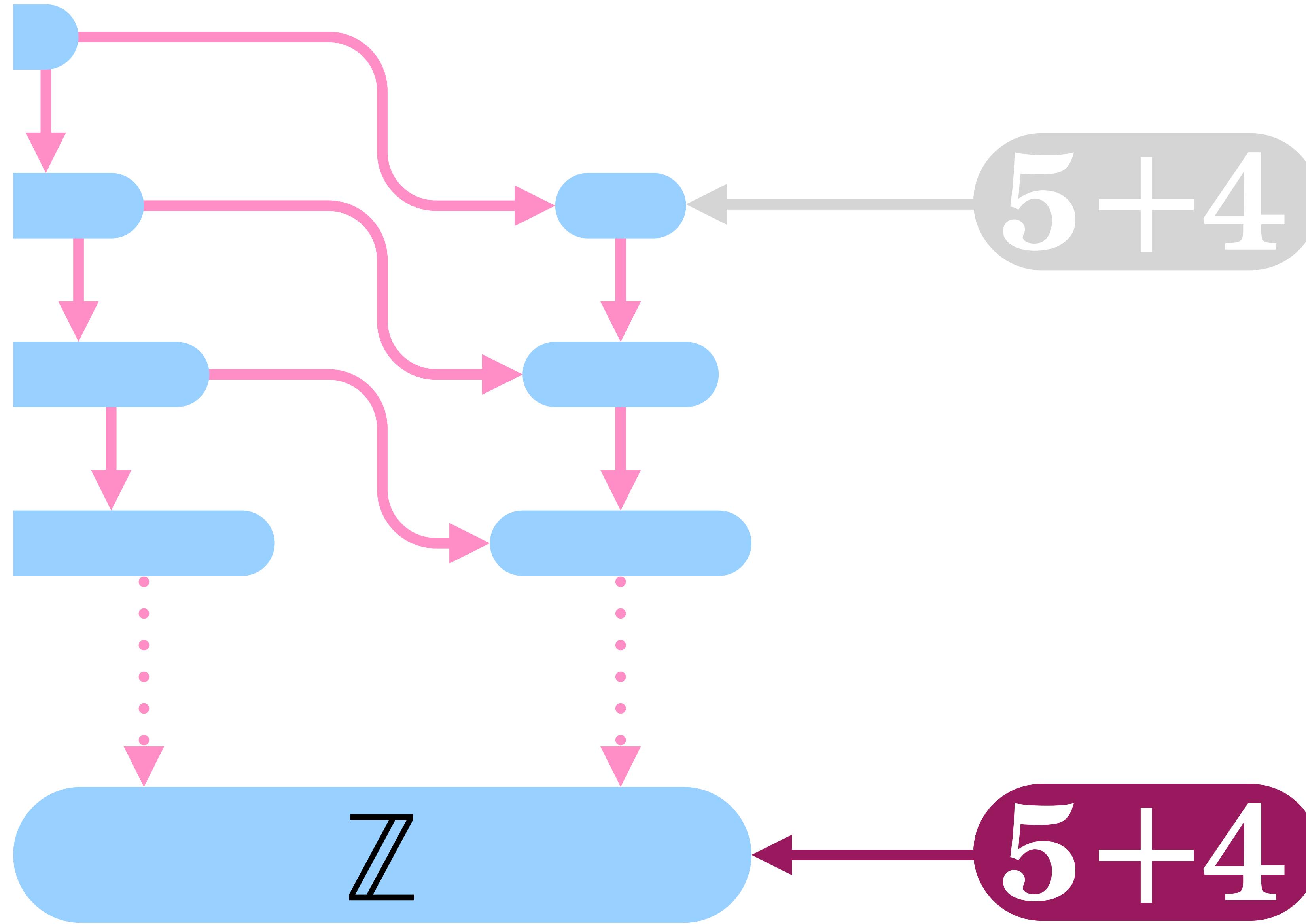


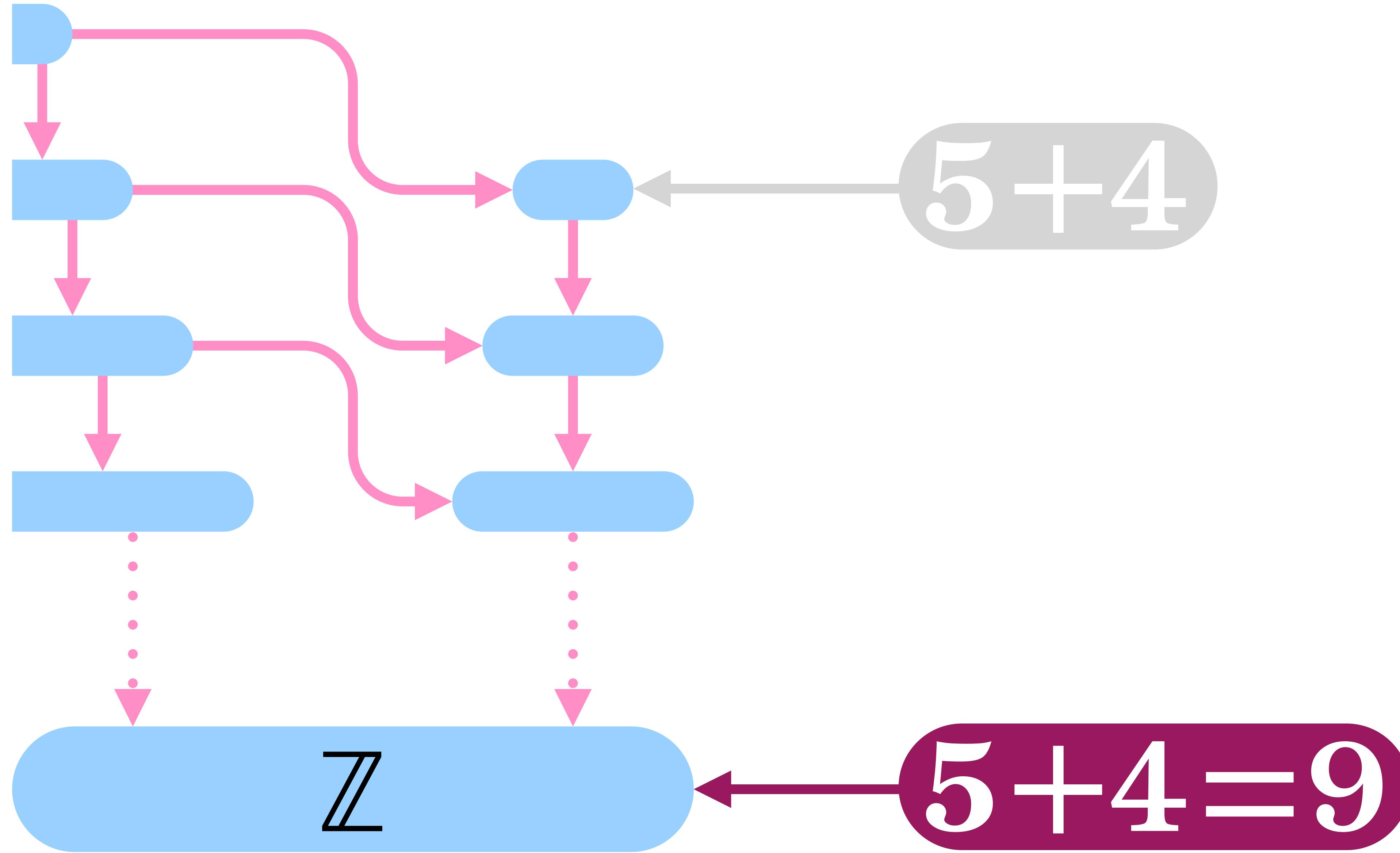


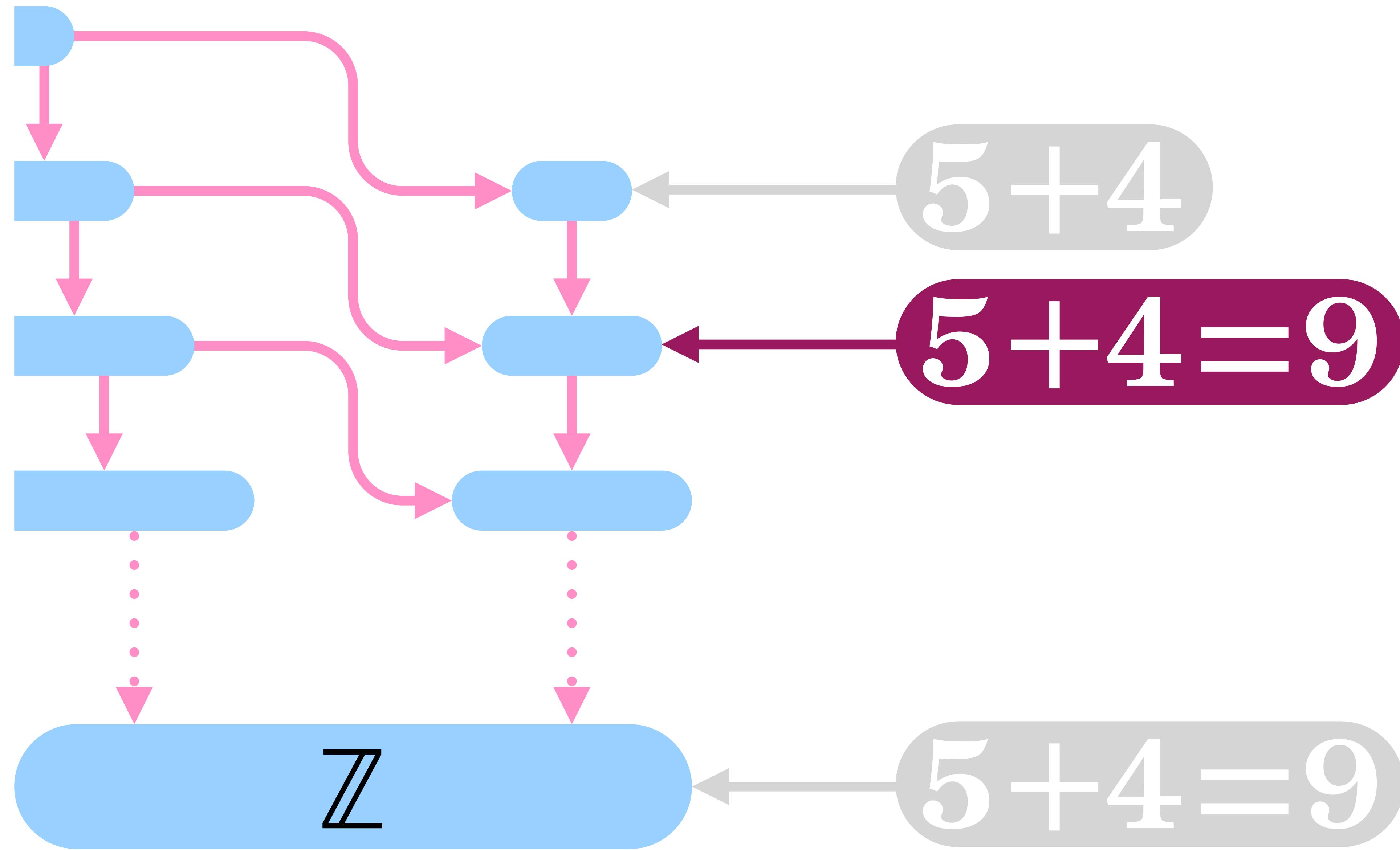


```
const auto result_in_Z = convert<Z>( a ) + convert<Z>( b );  
  
const auto expected_result = integer_kind_of<T>.is_signed()  
    ? convert_narrowing< T >( result_in_Z )  
    : convert_modular < T >( result_in_Z );  
  
// ...implementation...  
  
claim result == expected_result;
```









```
template <integer_kind k> class widening;
```

The types **widening<k>** do not have:

- 🚫 implicit narrowing
- 🚫 implicit modular arithmetic
- 🚫 modular conversion on assignment
- 🚫 integral promotion
- 🚫 usual arithmetic conversions
- 🚫 compound assignment operators
- 🚫 increment or decrement operators
- 🚫 bit fields
- 🚫 character types

The types **widening<k>** do have:

- ✓ integer kind k
- ✓ operations with results wider than their parameters
- ✓ operations that agree with \mathbb{Z}

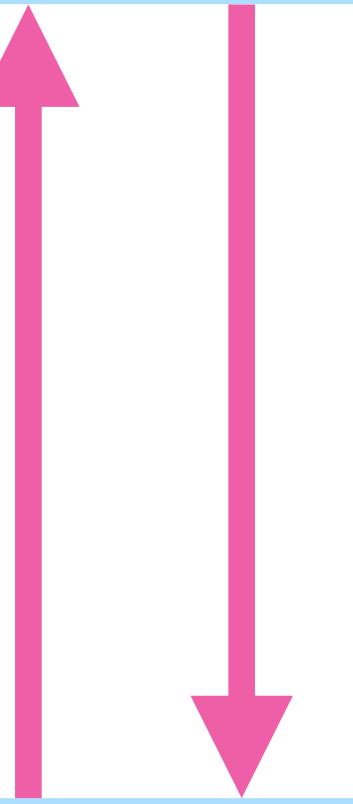
```
inline constexpr operator+( integer_kind ak, integer_kind bk )
{
    return join( ak, bk ) + one_bit;
}
```

```
template < integer_kind ak, integer_kind bk >
widenings<ak+bk> operator+( widenings<ak>& a, widenings<bk>& b );
```

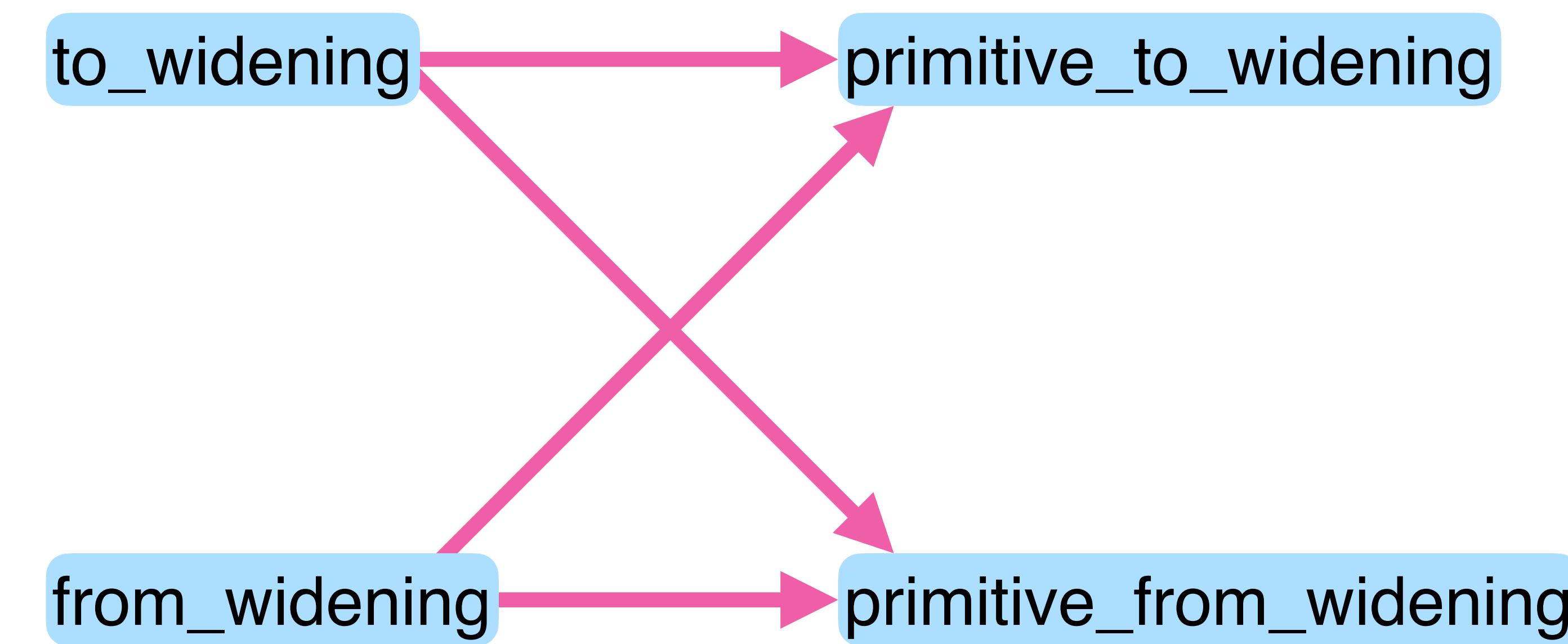
```
template < std::integral T >
widening< integral_kind_of<T> > to_widening( const T& a )
interface
{
    // ...implementation...
    claim from_widening< T >( result ) == a;
}
```

```
template < std::integral T >
T from_widening( const widening< integral_kind_of<T> >& a )
interface
{
    // ...implementation...
    claim to_widening( result ) == a;
}
```

to_widening



from_widening



```
template < std::integral T >
widening< integral_kind_of<T> > to_widening( const T& a )
interface
{
    // ...implementation...
    claim result == primitive_to_widening( a );
    claim primitive_from_widening< T >( result ) == a;
}
```

```
template < std::integral T >
T from_widening( const widening< integral_kind_of<T> >& a )
interface
{
    // ...implementation...
    claim result == primitive_from_widening< T >( a );
    claim primitive_to_widening( result ) == a;
}
```

```
const auto result_in_Z = convert<Z>( a ) + convert<Z>( b );  
  
const auto expected_result = integer_kind_of<T>.is_signed()  
    ? convert_narrowing< T >( result_in_Z )  
    : convert_modular < T >( result_in_Z );  
  
// ...implementation...  
  
claim result == expected_result;
```

```
const auto wide_result      = to_widening( a ) + to_widening( b );  
  
using lifted_type           = widening< integer_kind_of<T> >;  
  
const auto lifted_result    = integer_kind_of< T >.is_signed()  
    ? convert_narrowing< lifted_type >( wide_result )  
    : convert_modular < lifted_type >( wide_result );  
  
const auto expected_result = from_widening< T >( lifted_result );  
  
// ...implementation...  
  
claim result == expected_result;
```

- b	a + b	a & b	a == b
~b	a - b	a b	a < b
	a * b	a ^ b	a > b
next	a / b		a <= b
prev	a % b		a >= b
			a != b
convert			a <=> b
convert_modular			
convert_narrowing			

✖ Widening parameters

a << b
a >> b

✖ Widening × bit_size_t

The result of the binary + operator is the sum of the operands.

7.6.6 [expr.add]

```
template < integer_kind ak, integer_kind_bk >
widenings< ak + bk > operator+( const widenings<ak>& a,
                                    const widenings<bk>& b )
```

interface

```
{  
    claim_immunity a;                                // immunity from instability  
    claim_immunity b;                                // immunity from instability  
    discern a;                                       // discernible input  
    discern b;                                       // discernible input
```

```
const auto expected_result = add_with_carry( a, b );
```

implementation;

```
claim_right result;                                // right of stability  
discern result;                                   // discernible output
```

```
claim result == expected_result;                  // substitutability  
}
```



```
template < integer_kind ak, integer_kind bk >
widenings< ak + bk >  add_with_carry( const widenings<ak>& a,
                                         const widenings<bk>& b,
                                         const positive_bit& carry = {} )
```

interface

```
{
```

claim_immunity a;	// immunity from instability
claim_immunity b;	// immunity from instability
claim_immunity carry;	// immunity from instability
discern a;	// discernible input
discern b;	// discernible input
discern carry;	// discernible input

implementation;

claim_right result;	// right of stability
discern result;	// discernible output

// For a postcondition about the value of the result,
// call reference_add_axiom(a, b, carry);

```
}
```

```
template < integer_kind ak, integer_kind_bk >
void reference_add_axiom( const widening<ak>& a,
                           const widening<bk>& b,
                           const positive_bit& carry = {} )
```

interface

```
{
```

```
const auto reference_sum = reference_add_with_carry( a, b, carry );
```

The calling function is responsible
for the top part of the interface.

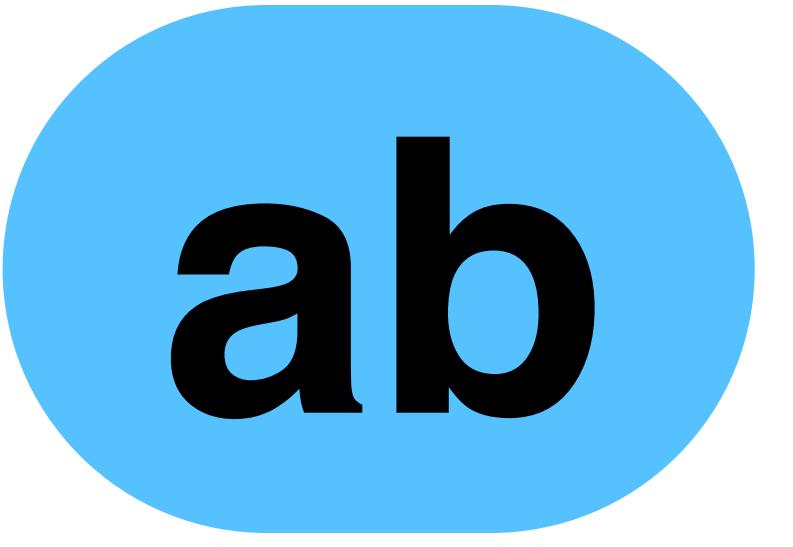
posit implementation:

There may be no function responsible
for the bottom part of the interface.

```
const auto sum = add_with_carry( a, b, carry );
```

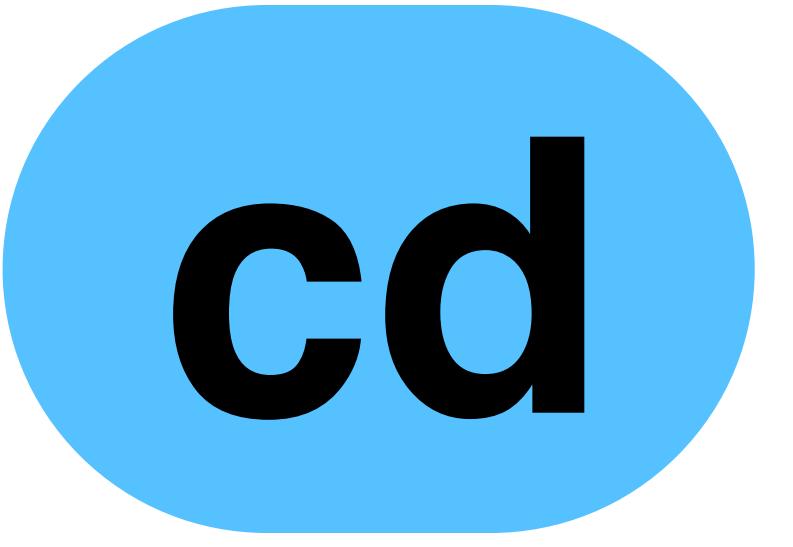
```
claim substitutable( sum, reference_sum );
```

```
}
```



ab

+



cd

0

ab

+

cd

0

ab

cd

a

c

0

b

d

0
ab
cd

a
c

0
b
d

+

18

0
ab
cd

a
c

0
b
d
+
—

1 8

0
ab
cd

1
a
c

0
b
d

+

1
8

0
ab
cd

$$\begin{array}{r} 1 \\ a \\ c \\ \hline 17 \end{array}$$

$$\begin{array}{r} 0 \\ b \\ d \\ \hline 1 \\ 8 \end{array}$$

0

ab

+ cd

178

1

a

+ c

17

0

b

+ d

1 8

```
template < integer_kind abk, integer_kind cdk >
inline auto reference_add_with_carry( const widening<abk>& ab,
                                      const widening<cdk>& cd,
                                      const positive_bit& x0 = {} )
{
    if constexpr ( abk.width() > 1 && cdk.width() > 1 && lo_width( abk ) == lo_width( cdk ) )
    {
        const auto [ a, b ] = split_bits( ab );
        const auto [ c, d ] = split_bits( cd );

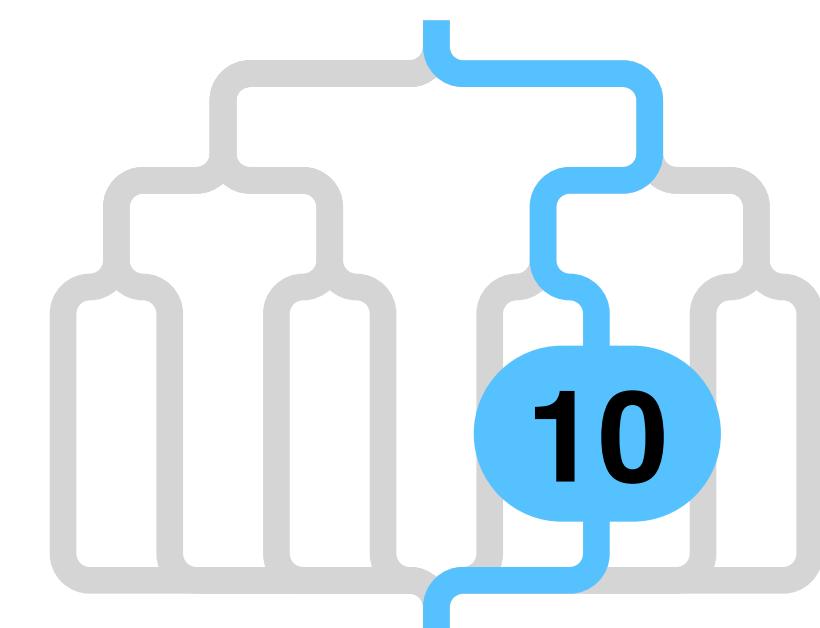
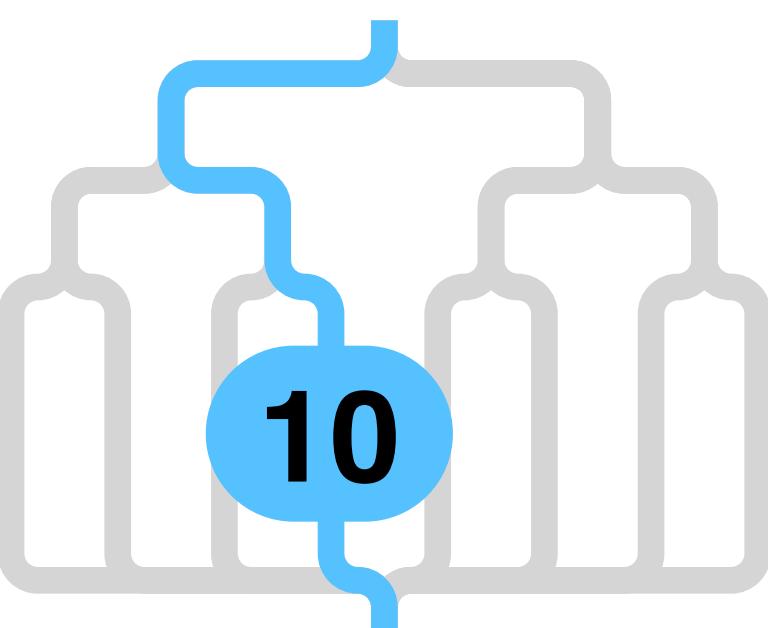
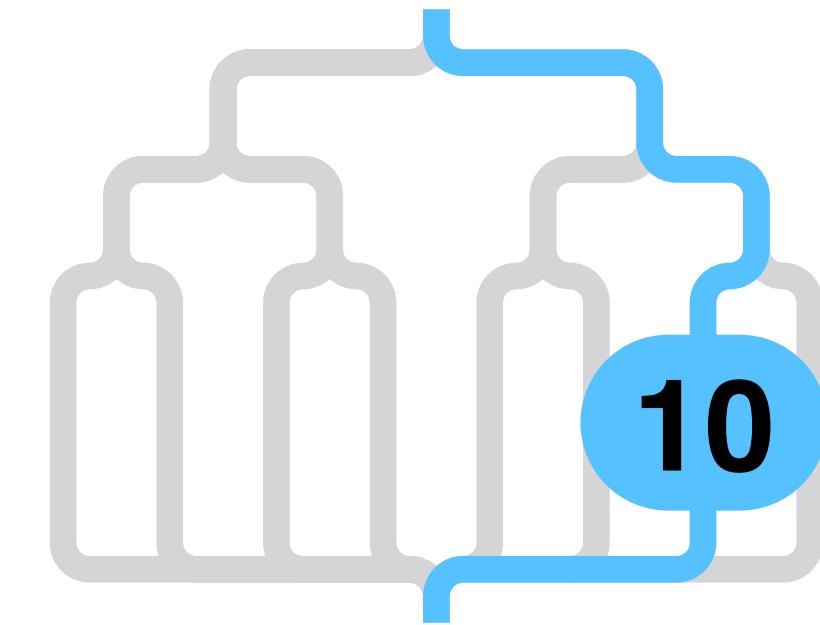
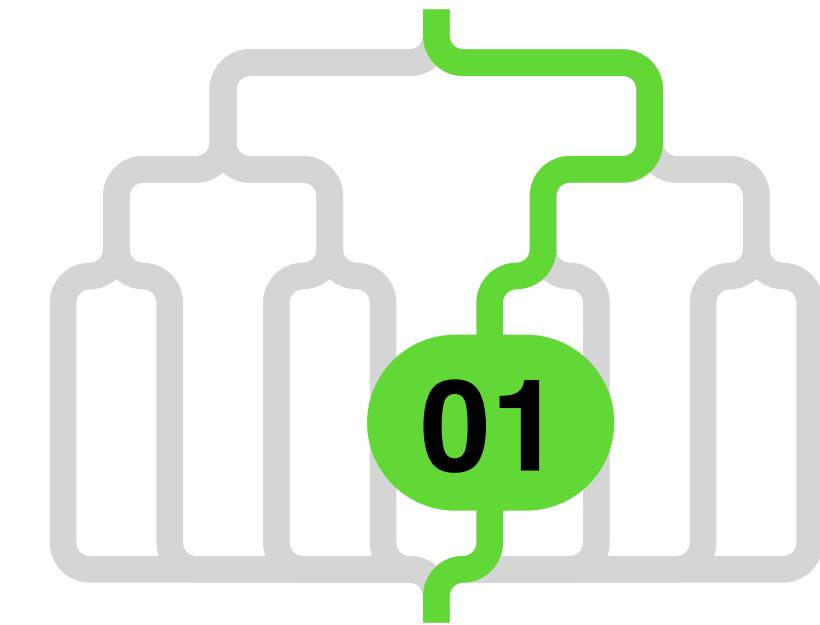
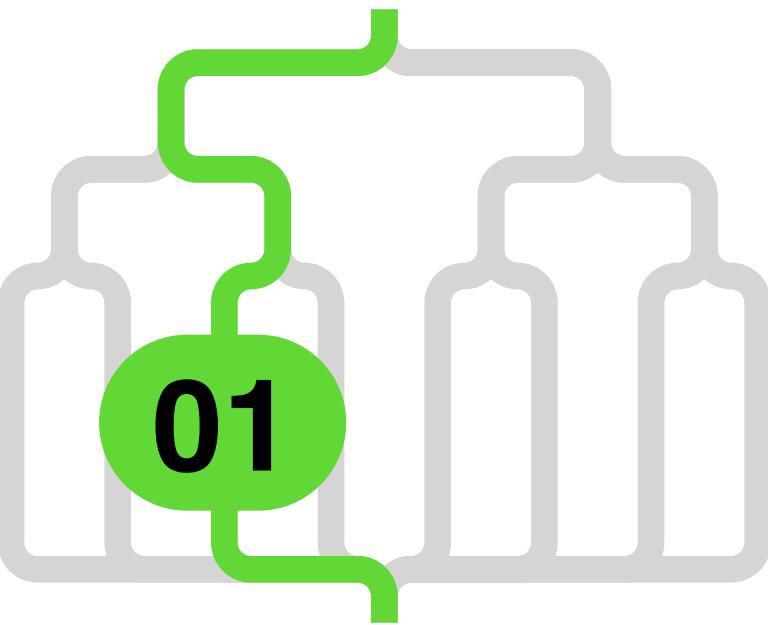
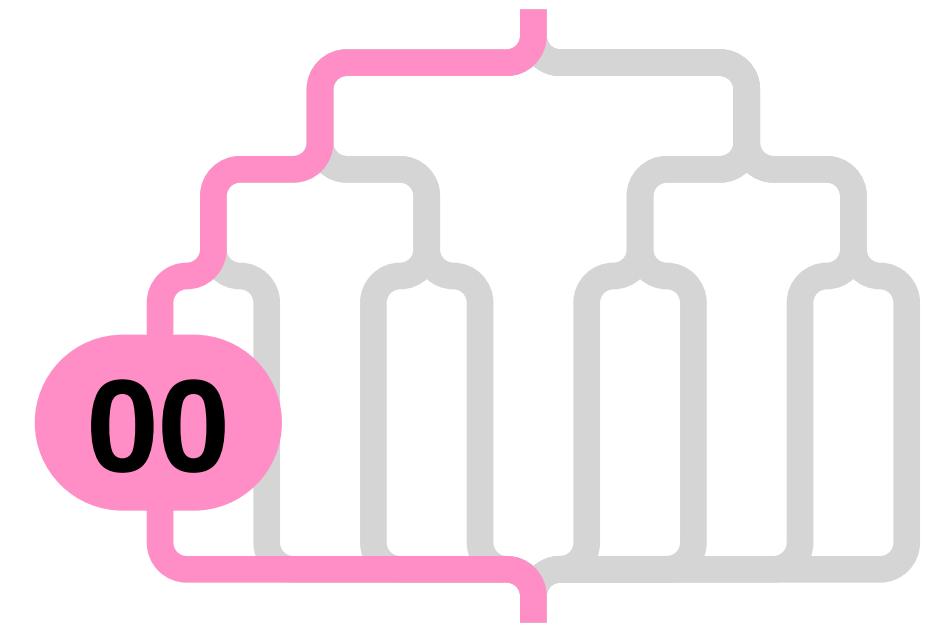
        const auto s0      = add_with_carry( b, d, x0 );
        const auto [ x1, r0 ] = split_bits( s0 );

        const auto r1      = add_with_carry( a, c, x1 );

        return join_bits( r1, r0 );
    }
    else
        // ...
}
```

```
//...
else if constexpr ( abk == cdk && abk.width() == one_bit && !abk.is_signed() )
{
    const auto bit_0 = bool_to_bit( false );
    const auto bit_1 = bool_to_bit( true );

    if ( bit_to_bool( ab ) )
        if ( bit_to_bool( cd ) )
            if ( bit_to_bool( x0 ) )
                return join_bits( bit_1, bit_1 );
        else
            return join_bits( bit_1, bit_0 );
    else
        if ( bit_to_bool( x0 ) )
            return join_bits( bit_1, bit_0 );
        else
            return join_bits( bit_0, bit_1 );
}
else
    if ( bit_to_bool( cd ) )
        if ( bit_to_bool( x0 ) )
            return join_bits( bit_1, bit_0 );
        else
            return join_bits( bit_0, bit_1 );
    else
        if ( bit_to_bool( x0 ) )
            return join_bits( bit_0, bit_1 );
        else
            return join_bits( bit_0, bit_0 );
//...
```





This is all that's left

split_bits
join_bits

bit_to_bool
bool_to_bit

sign_bit_to_bool
bool_to_sign_bit

```
template < integer_kind ak, integer_kind bk >
void addition_is_commutative( const widening<ak>& a,
                             const widening<bk>& b,
                             const positive_bit& x = {} )
```

interface

```
{  
    const auto a_plus_b = add_with_carry( a, b, x );  
    const auto b_plus_a = add_with_carry( b, a, x );
```

The calling function is responsible
for the top part of the interface.

claim implementation;

The called function is responsible
for the bottom part of the interface.

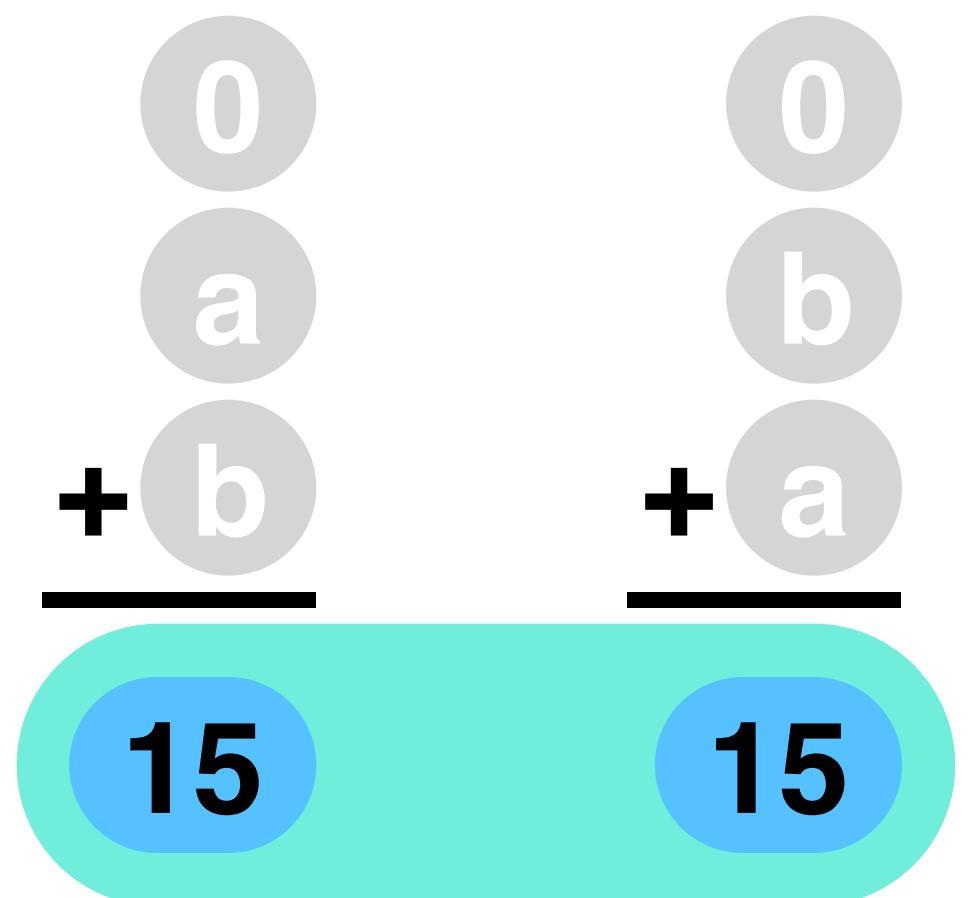
```
claim substitutable( a_plus_b, b_plus_a );  
}
```

```
template < integer_kind ak, integer_kind bk >
void addition_is_commutative( const widening<ak>& a,
                             const widening<bk>& b,
                             const positive_bit& x = {} )
```

interface

```
{
```

```
const auto a_plus_b = add_with_carry( a, b, x );
const auto b_plus_a = add_with_carry( b, a, x );
```



claim implementation;

```
claim substitutable( a_plus_b, b_plus_a );
}
```

0
ab
cd

```
reference_add_axiom( ab, cd, 0 );
```

$$\begin{array}{r} 0 \\ ab \\ + cd \\ \hline 178 \end{array}$$

$$\begin{array}{r} 1 \\ a \\ + c \\ \hline 17 \end{array}$$

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline 18 \end{array}$$

```
reference_add_axiom( ab, cd, 0 );
```

$$\begin{array}{r} 0 \\ ab \\ + cd \\ \hline 178 \end{array}$$

$$\begin{array}{r} 1 \\ a \\ + c \\ \hline 17 \end{array}$$

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline 18 \end{array}$$

$$\begin{array}{r} 0 \\ ab \\ + cd \\ \hline \end{array}$$

$$\begin{array}{r} a \\ + c \\ \hline \end{array}$$

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline \end{array}$$

```
reference_add_axiom( ab, cd, 0 );
```

$$\begin{array}{r}
 0 \\
 ab \\
 + cd \\
 \hline
 178
 \end{array}
 \quad
 \begin{array}{r}
 1 \\
 a \\
 + c \\
 \hline
 17
 \end{array}
 \quad
 \begin{array}{r}
 0 \\
 b \\
 + d \\
 \hline
 18
 \end{array}$$

```
reference_add_axiom( ab, cd, 0 );
```

$$\begin{array}{r}
 0 \\
 b \\
 + d \\
 \hline
 18
 \end{array}
 \quad
 \begin{array}{r}
 0 \\
 d \\
 + b \\
 \hline
 18
 \end{array}$$

```
addition_is_commutative( b, d, 0 );
```

$$\begin{array}{r}
 0 \\
 ab \\
 + cd \\
 \hline
 a \\
 + c \\
 \hline
 b \\
 + d \\
 \hline
 0
 \end{array}$$

$$\begin{array}{r} 0 \\ ab \\ + cd \\ \hline 178 \end{array}$$

$$\begin{array}{r} 1 \\ a \\ + c \\ \hline 17 \end{array}$$

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline 18 \end{array}$$

$$\begin{array}{r} 0 \\ ab \\ cd \\ \hline \end{array}$$

$$\begin{array}{r} 1 \\ a \\ c \\ \hline \end{array}$$

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline 18 \end{array}$$

```
reference_add_axiom( ab, cd, 0 );
```

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline 18 \end{array}$$

$$\begin{array}{r} 0 \\ d \\ + b \\ \hline 18 \end{array}$$

```
addition_is_commutative( b, d, 0 );
```

$$\begin{array}{r}
 0 \\
 ab \\
 + cd \\
 \hline
 178
 \end{array}
 \quad
 \begin{array}{r}
 1 \\
 a \\
 + c \\
 \hline
 17
 \end{array}
 \quad
 \begin{array}{r}
 0 \\
 b \\
 + d \\
 \hline
 18
 \end{array}$$

`reference_add_axiom(ab, cd, 0);`

$$\begin{array}{r}
 0 \\
 b \\
 + d \\
 \hline
 18
 \end{array}
 \quad
 \begin{array}{r}
 0 \\
 d \\
 + b \\
 \hline
 18
 \end{array}$$

$$\begin{array}{r}
 1 \\
 a \\
 + c \\
 \hline
 17
 \end{array}
 \quad
 \begin{array}{r}
 1 \\
 c \\
 + a \\
 \hline
 17
 \end{array}$$

`addition_is_commutative(b, d, 0);`

`addition_is_commutative(a, c, 1);`

$$\begin{array}{r} 0 \\ ab \\ + cd \\ \hline 178 \end{array}$$

$$\begin{array}{r} 1 \\ a \\ + c \\ \hline 17 \end{array}$$

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline 18 \end{array}$$

$$\begin{array}{r} 0 \\ cd \\ + ab \\ \hline 178 \end{array}$$

$$\begin{array}{r} 1 \\ c \\ + a \\ \hline 17 \end{array}$$

$$\begin{array}{r} 0 \\ d \\ + b \\ \hline 18 \end{array}$$

reference_add_axiom(ab, cd, 0);

reference_add_axiom(cd, ab, 0);

$$\begin{array}{r} 0 \\ b \\ + d \\ \hline 18 \end{array}$$

$$\begin{array}{r} 0 \\ d \\ + b \\ \hline 18 \end{array}$$

$$\begin{array}{r} 1 \\ a \\ + c \\ \hline 17 \end{array}$$

$$\begin{array}{r} 1 \\ c \\ + a \\ \hline 17 \end{array}$$

addition_is_commutative(b, d, 0);

addition_is_commutative(a, c, 1);

```
template < integer_kind abk, integer_kind cdk >
void addition_is_commutative( const widening<abk>& ab, const widening<cdk>& cd,
                             const positive_bit& x0 = {} )

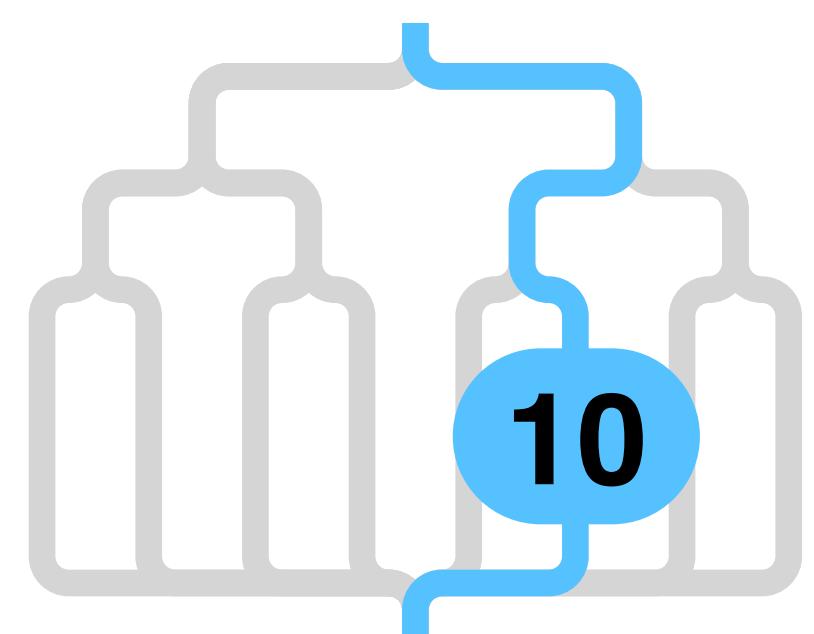
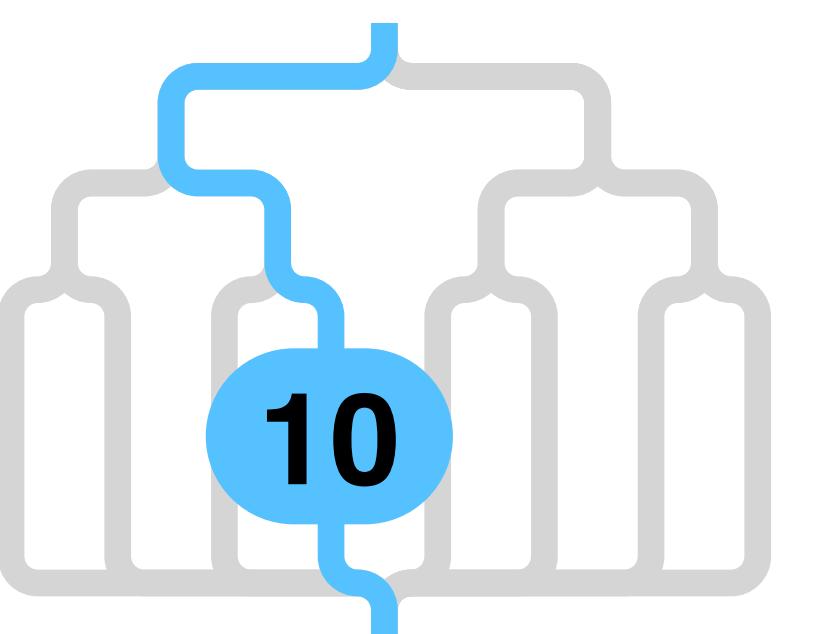
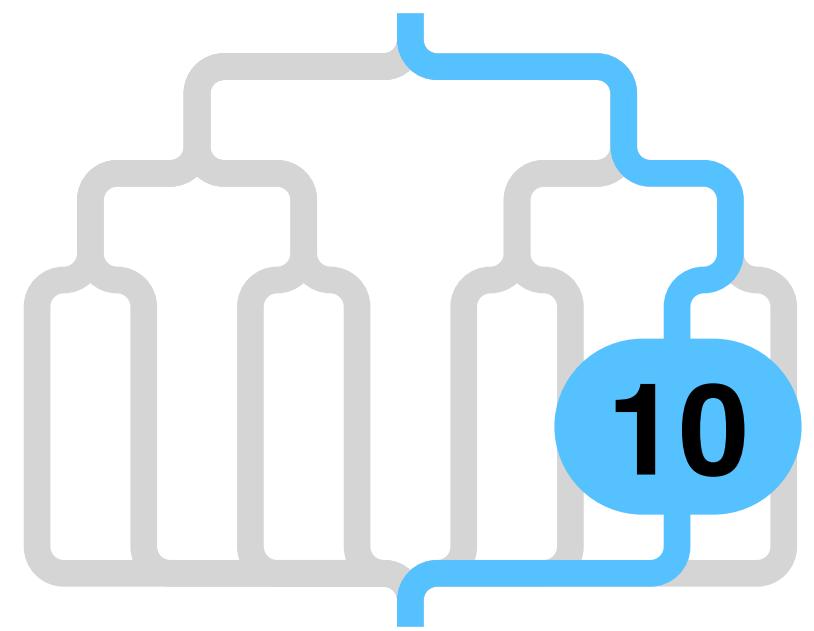
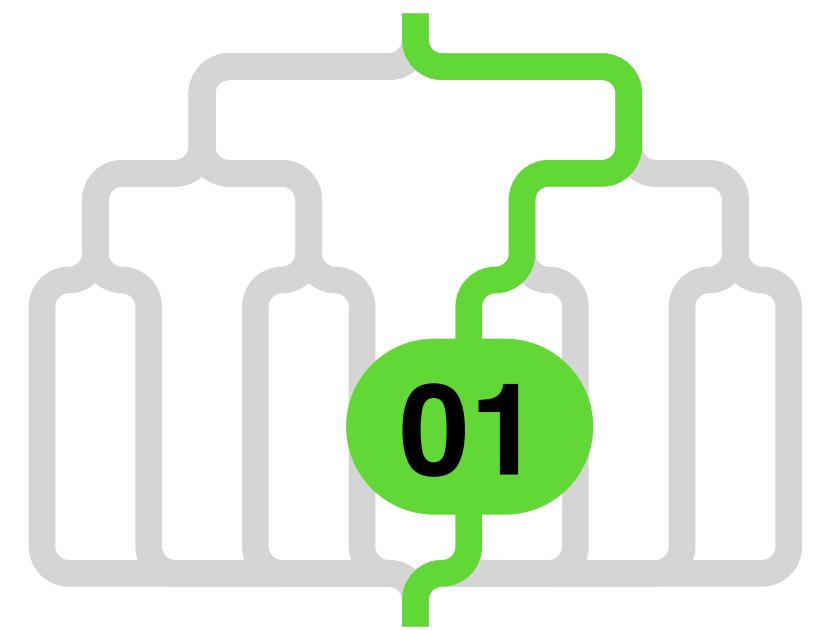
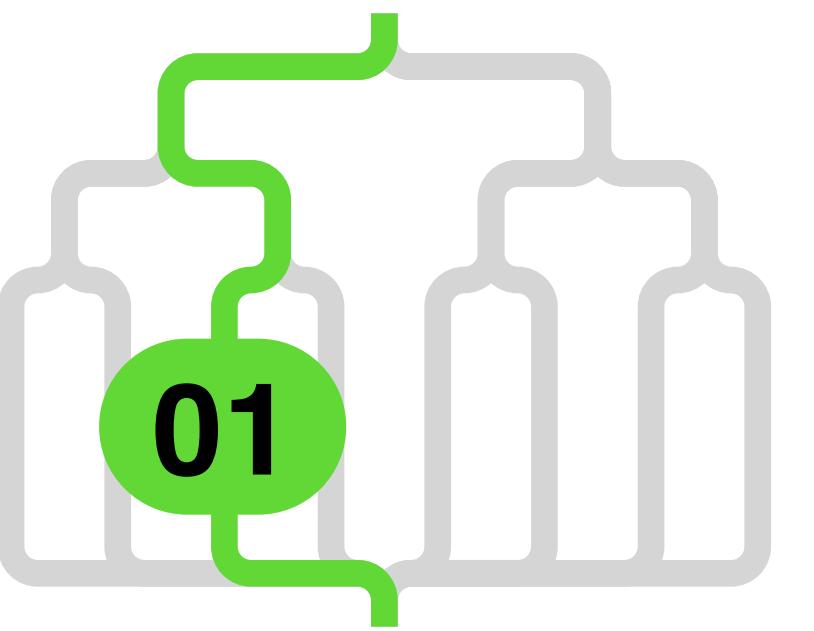
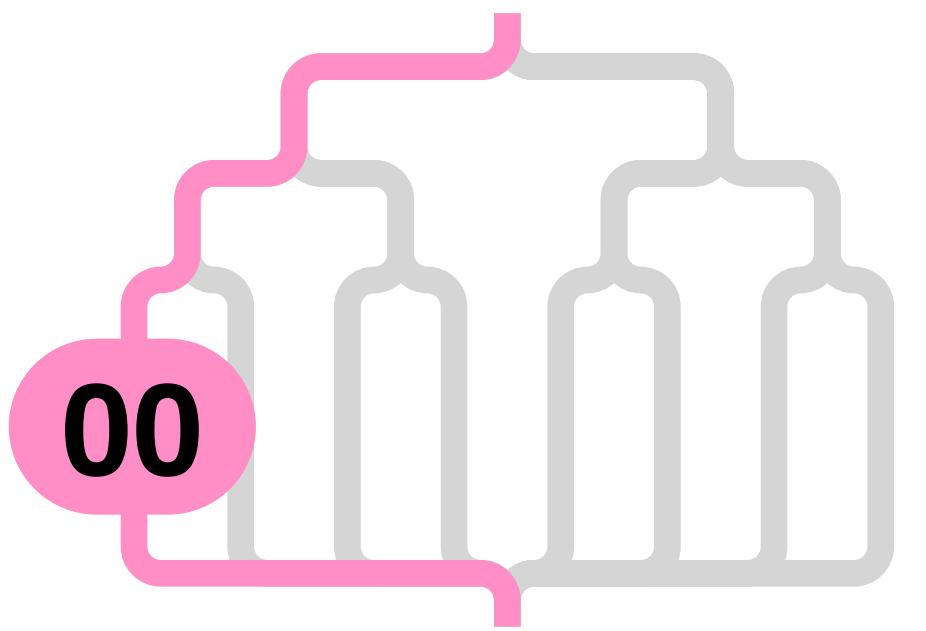
implementation

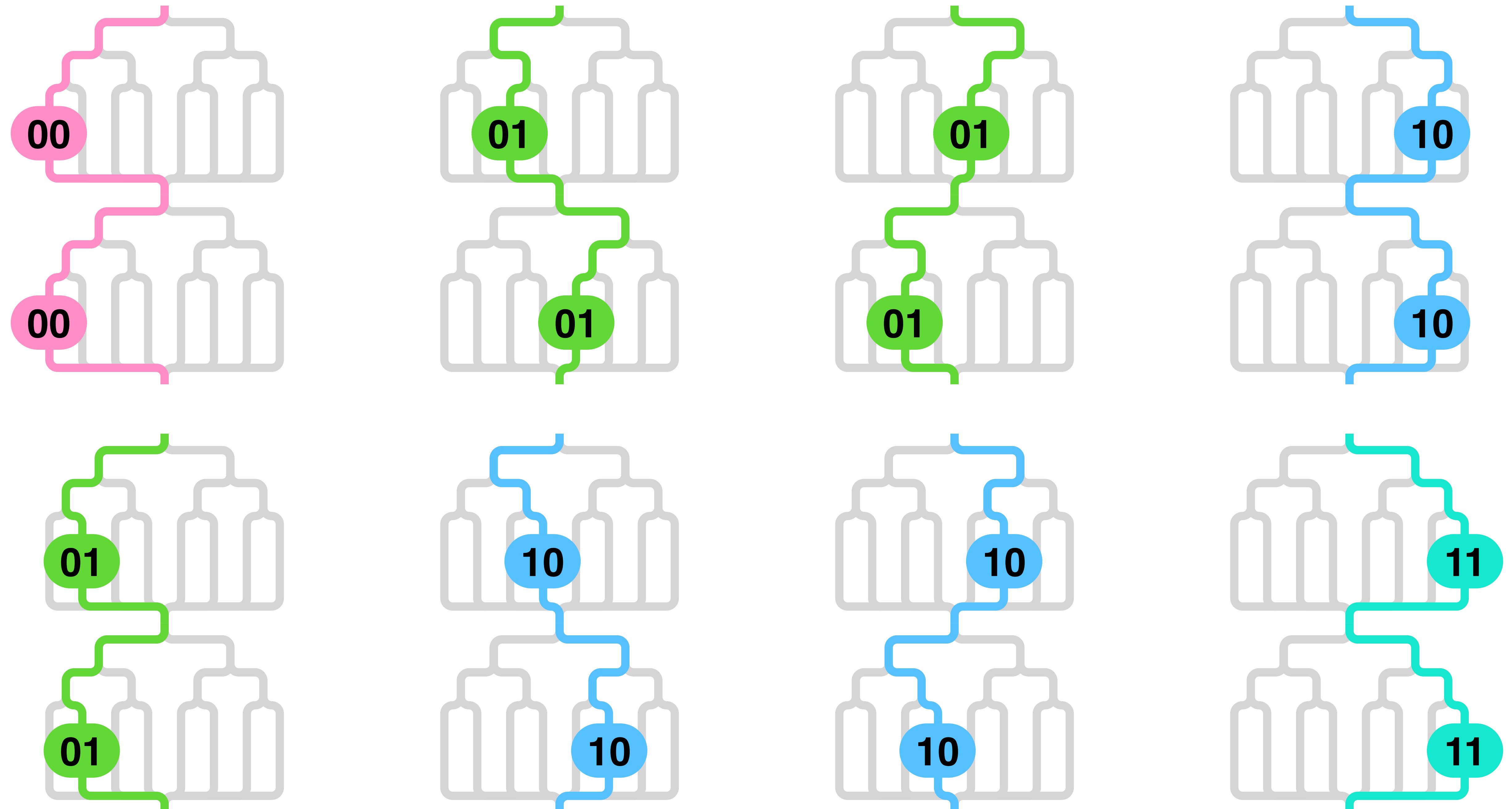
{
    if constexpr ( abk.width() > 1 && cdk.width() > 1 && lo_width( abk ) == lo_width( cdk ) )
    {
        reference_add_axiom( ab, cd, x0 );

        const auto [ a, b ] = split_bits( ab );
        const auto [ c, d ] = split_bits( cd );
        addition_is_commutative( b, d, x0 );

        const auto s0      = add_with_carry( b, d, x0 );
        const auto [ x1, r0 ] = split_bits( s0 );
        addition_is_commutative( a, c, x1 );

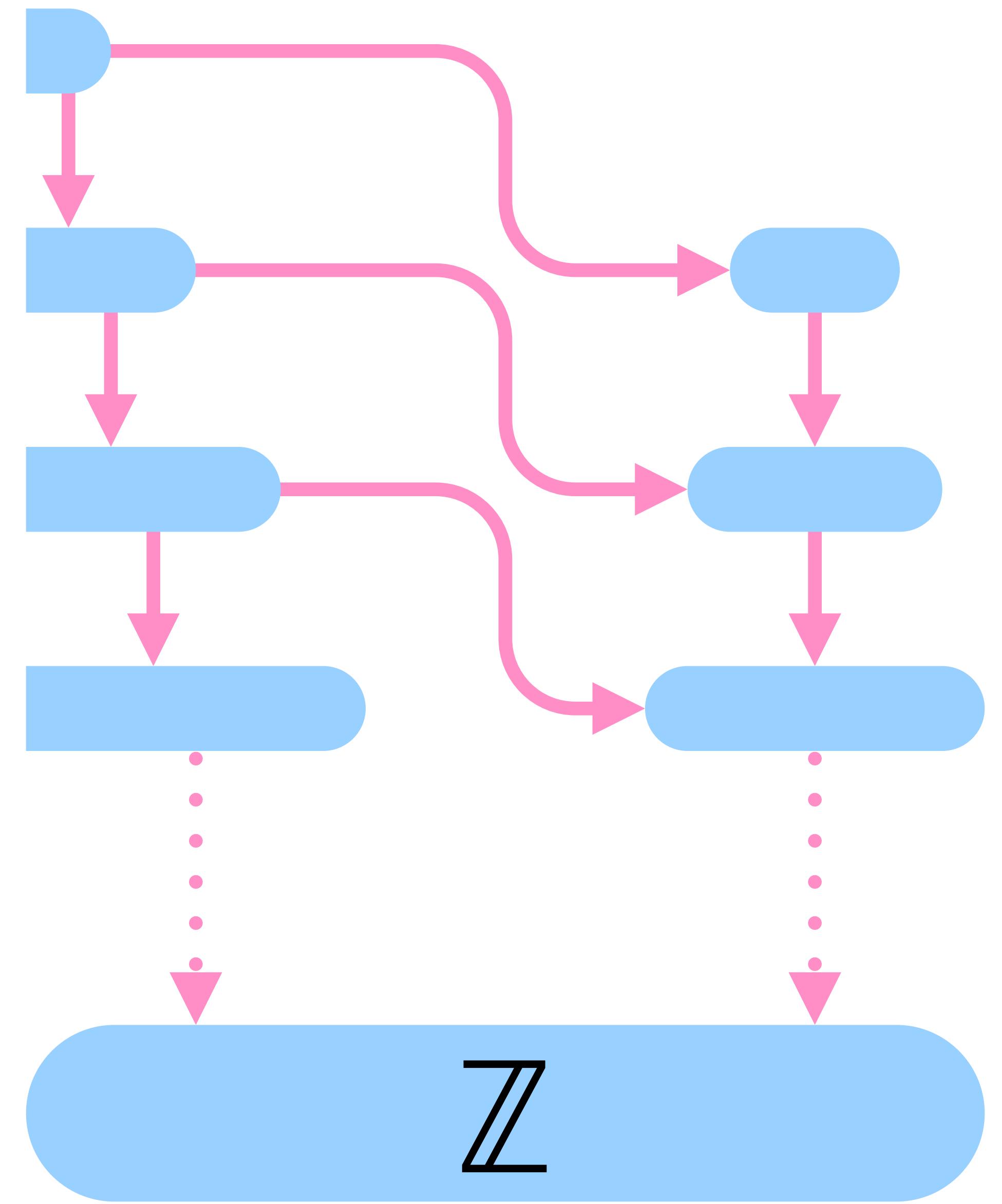
        reference_add_axiom( cd, ab, x0 );
    }
    else // ...
}
```

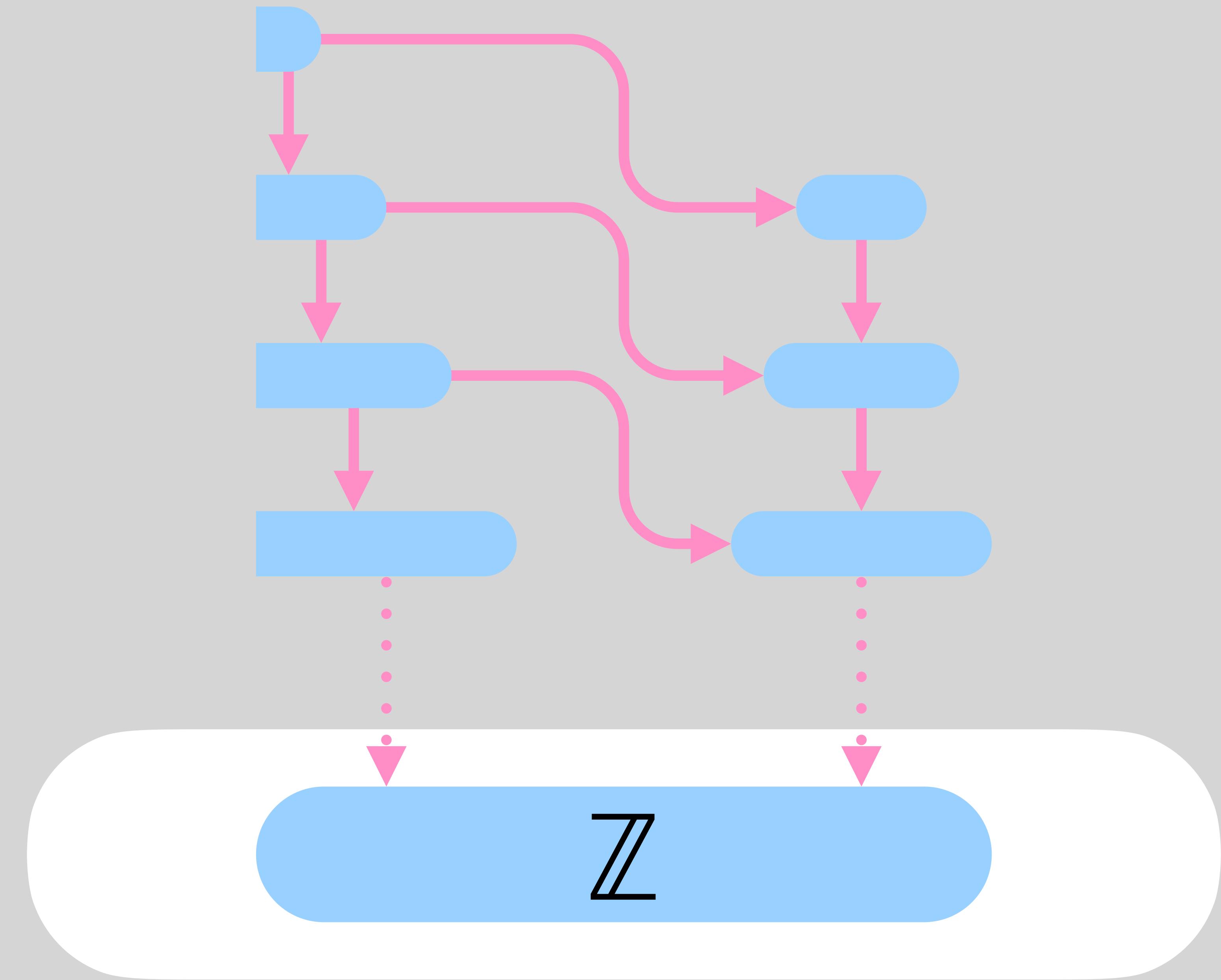


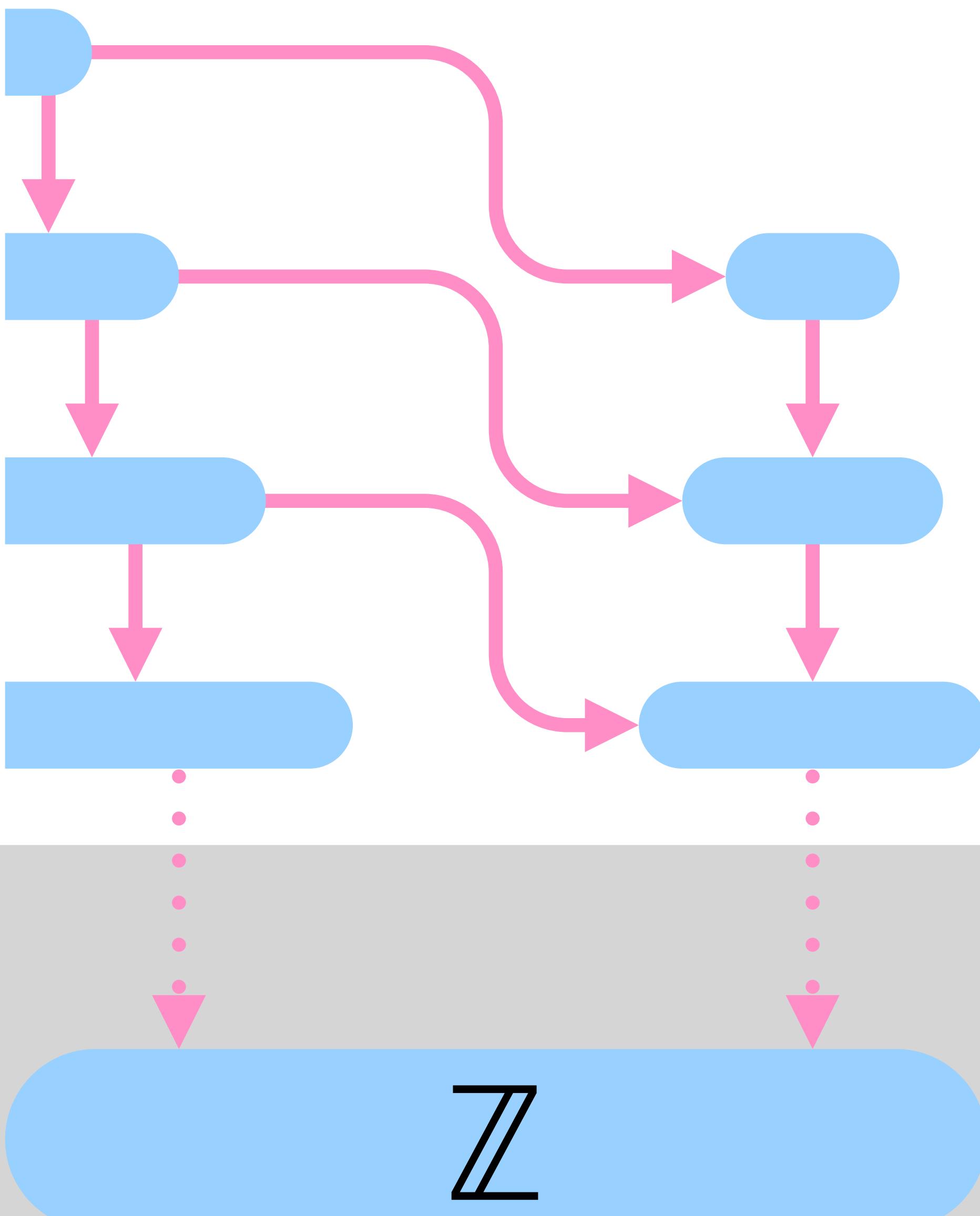


```
// ...
else if constexpr ( abk == cdk && abk.width() == one_bit && !abk.is_signed() )
{
    reference_add_axiom( ab, cd, x0 );

    reference_add_axiom( cd, ab, x0 );
}
else // ...
```







Thank you for listening.

Questions?